

Relations

Relations

We use sets of pairs called relations to model relationships between objects.

The presence of a pair of objects in the set indicates that the first is related to the second.

Notation

If X and Y are sets, then $X \leftrightarrow Y$ denotes the set of all relations between X and Y :

$$X \leftrightarrow Y == \mathbb{P}(X \times Y)$$

The pair (p, q) can also be written as $p \mapsto q$.

Example

$$\begin{aligned}
 \{a, b\} \leftrightarrow \{0, 1\} = & \{\emptyset, \\
 & \{(a, 0)\}, \{(a, 1)\}, \{(b, 0)\}, \{(b, 1)\}, \\
 & \{(a, 0), (a, 1)\}, \{(a, 0), (b, 0)\}, \\
 & \{(a, 0), (b, 1)\}, \{(a, 1), (b, 0)\}, \\
 & \{(a, 1), (b, 1)\}, \{(b, 0), (b, 1)\}, \\
 & \{(a, 0), (a, 1), (b, 0)\}, \\
 & \{(a, 0), (a, 1), (b, 1)\}, \\
 & \{(a, 0), (b, 0), (b, 1)\}, \\
 & \{(a, 1), (b, 0), (b, 1)\}, \\
 & \{(a, 0), (a, 1), (b, 0), (b, 1)\}}
 \end{aligned}$$

Example

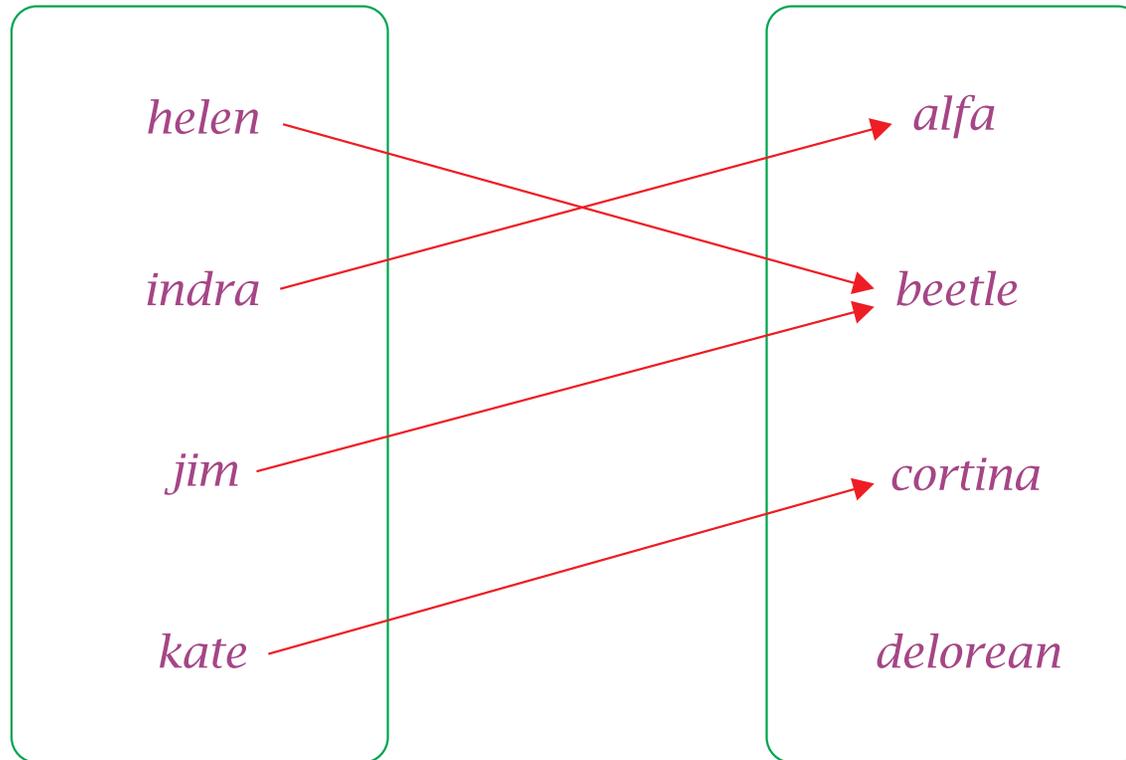
Drivers == { *helen*, *indra*, *jim*, *kate* }

Cars == { *alfa*, *beetle*, *cortina*, *delorean* }

drives == { *helen* \mapsto *beetle*, *indra* \mapsto *alfa*,
jim \mapsto *beetle*, *kate* \mapsto *cortina* }

Drivers

Cars



Domain and range

domain:

$$\text{dom } R = \{ x : X; y : Y \mid x \mapsto y \in R \bullet x \}$$

range:

$$\text{ran } R = \{ x : X; y : Y \mid x \mapsto y \in R \bullet y \}$$

Example

$\text{dom } drives = \{helen, indra, jim, kate\}$

$\text{ran } drives = \{alfa, beetle, cortina\}$

Restrictions

domain restriction:

$$A \triangleleft R = \{ x : X; y : Y \mid x \mapsto y \in R \wedge x \in A \bullet x \mapsto y \}$$

range restriction:

$$R \triangleright B = \{ x : X; y : Y \mid x \mapsto y \in R \wedge y \in B \bullet x \mapsto y \}$$

Subtraction

domain subtraction:

$$A \triangleleft R = \{ x : X; y : Y \mid x \mapsto y \in R \wedge x \notin A \bullet x \mapsto y \}$$

range subtraction:

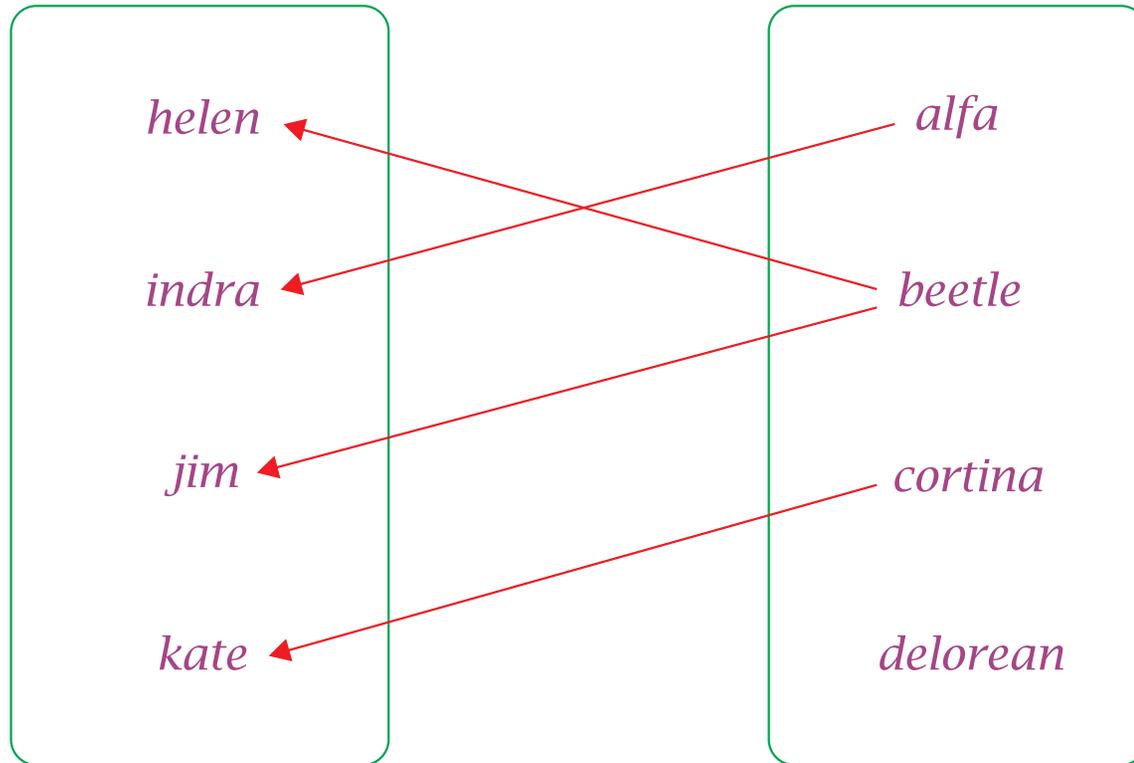
$$R \triangleright B = \{ x : X; y : Y \mid x \mapsto y \in R \wedge y \notin B \bullet x \mapsto y \}$$

Inverse

$$\forall x: X; y: Y \bullet x \mapsto y \in R^{\sim} \Rightarrow y \mapsto x \in R$$

Drivers

Cars



Relational image

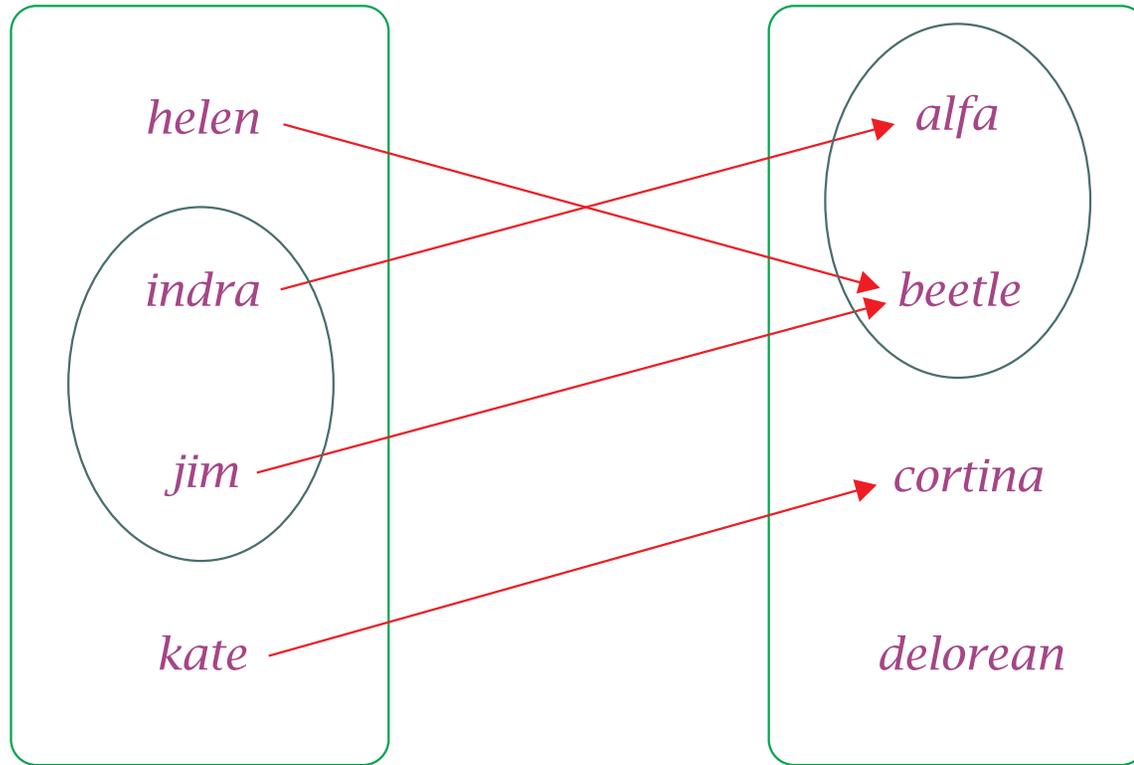
$$R(A) = \text{ran}(A \triangleleft R)$$

Example

$drives(\{indra, jim\}) = \{alfa, beetle\}$

Drivers

Cars



Question

What do the following sets look like in extension?

- $\{jim, kate\} \triangleleft drives$
- $drives \triangleright \{beetle\}$
- $\text{dom}(drives^{\sim})$
- $drives^{\sim} \setminus (\{beetle, delorean\})$

Relational composition

If the target of one relation matches the source of another, it may be useful to consider their relational composition.

The composition of two relations relates objects in the source of the first to objects in the target of the second, provided that some intermediate point exists.

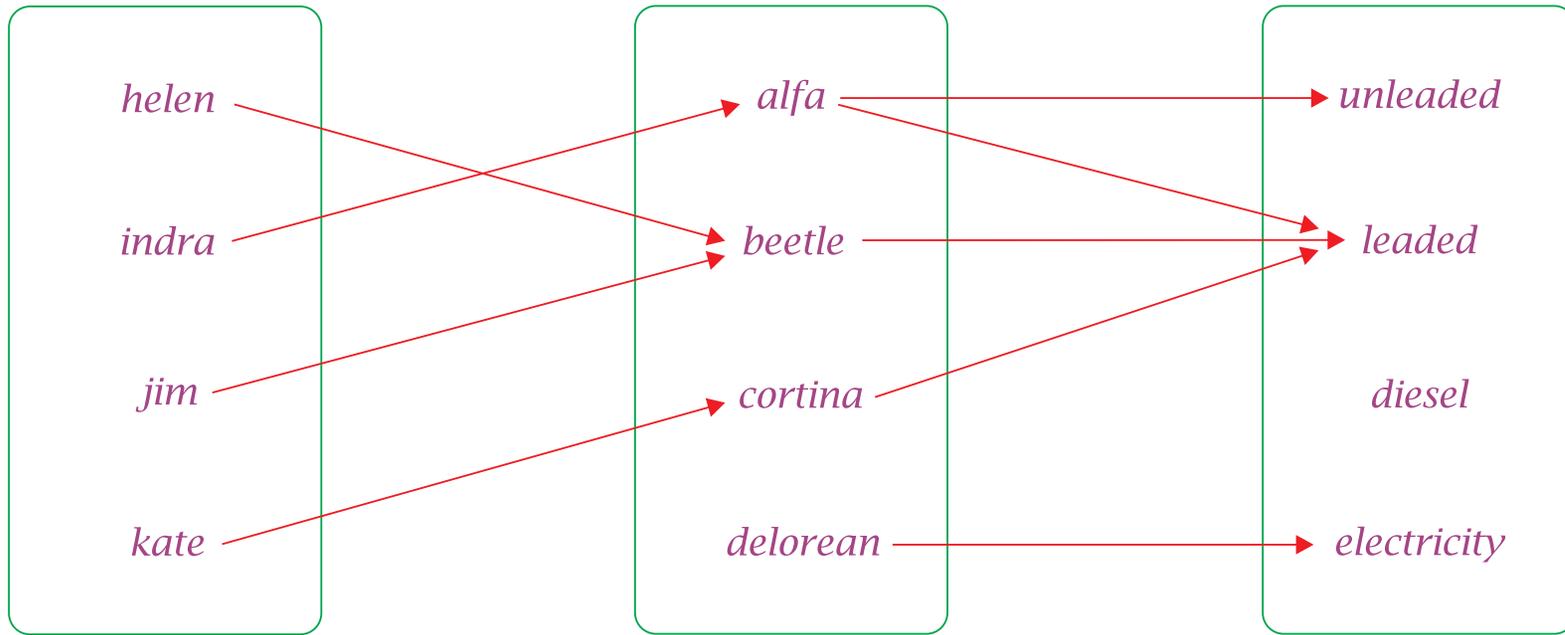
Definition

$$x \mapsto z \in R \circ S \Leftrightarrow \exists y: Y \bullet x \mapsto y \in R \wedge y \mapsto z \in S$$

Drivers

Cars

Fuels



drives

uses

Question

Which of the following statements are true?

- $jim \mapsto unleaded \in (drives \circ uses)$
- $electricity \in \text{ran}(drives \circ uses)$
- $indra \mapsto unleaded \in (drives \circ uses)$

Properties

- reflexivity
- symmetry (and ...)
- transitivity

Reflexivity and symmetry

$$\begin{aligned} \text{Reflexive}[X] == \{ R : X \leftrightarrow X \mid \\ \forall x : X \bullet \\ x \mapsto x \in R \} \end{aligned}$$

$$\begin{aligned} \text{Symmetric}[X] == \{ R : X \leftrightarrow X \mid \\ \forall x, y : X \bullet \\ x \mapsto y \in R \Rightarrow y \mapsto x \in R \} \end{aligned}$$

Transitivity

$$\text{Transitive}[X] == \{ R : X \leftrightarrow X \mid$$

$$\forall x, y, z : X \bullet$$

$$x \mapsto y \in R \wedge y \mapsto z \in R \Rightarrow$$

$$x \mapsto z \in R \}$$

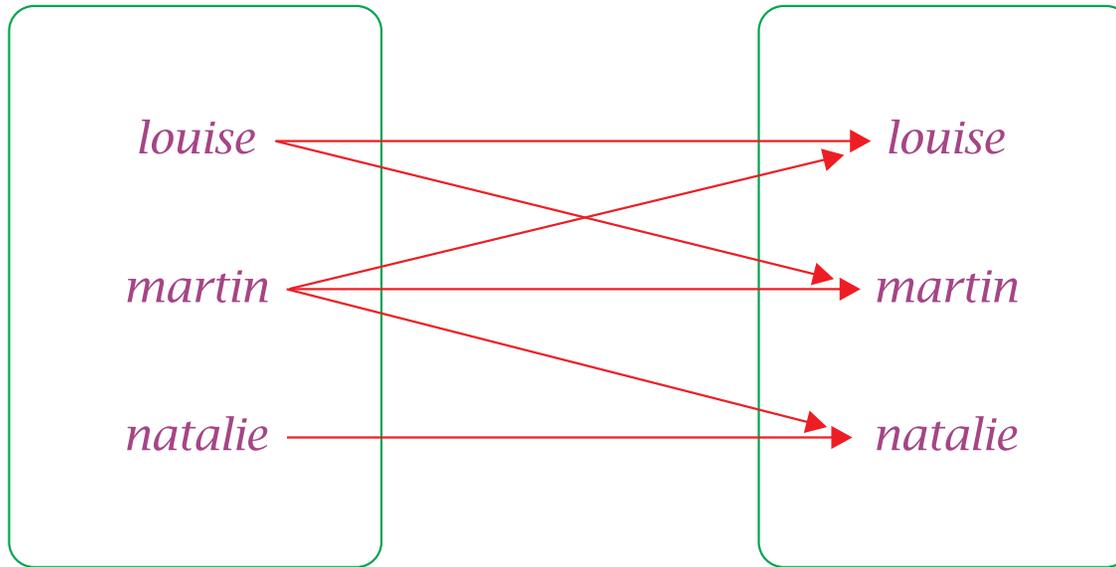
$$\text{Equivalence}[X] == \text{Reflexive}[X] \cap$$

$$\text{Symmetric}[X] \cap$$

$$\text{Transitive}[X]$$

House

House



Closure

If a relation does not have a specified property, then we may add maplets until it does.

If we add precisely those maplets that are lacking, then the resulting relation is called a **closure**.

Reflexive closure

To obtain the reflexive closure R^r of a relation R , we have only to add the maplets of the identity relation:

$$\text{id } X == \{ x : X \bullet x \mapsto x \}$$

Question

What should we add in order to obtain the symmetric closure of a relation S ?

Iteration

$$R^1 = R$$

$$R^2 = R \circ R$$

$$R^3 = R \circ R \circ R$$

$$\vdots$$

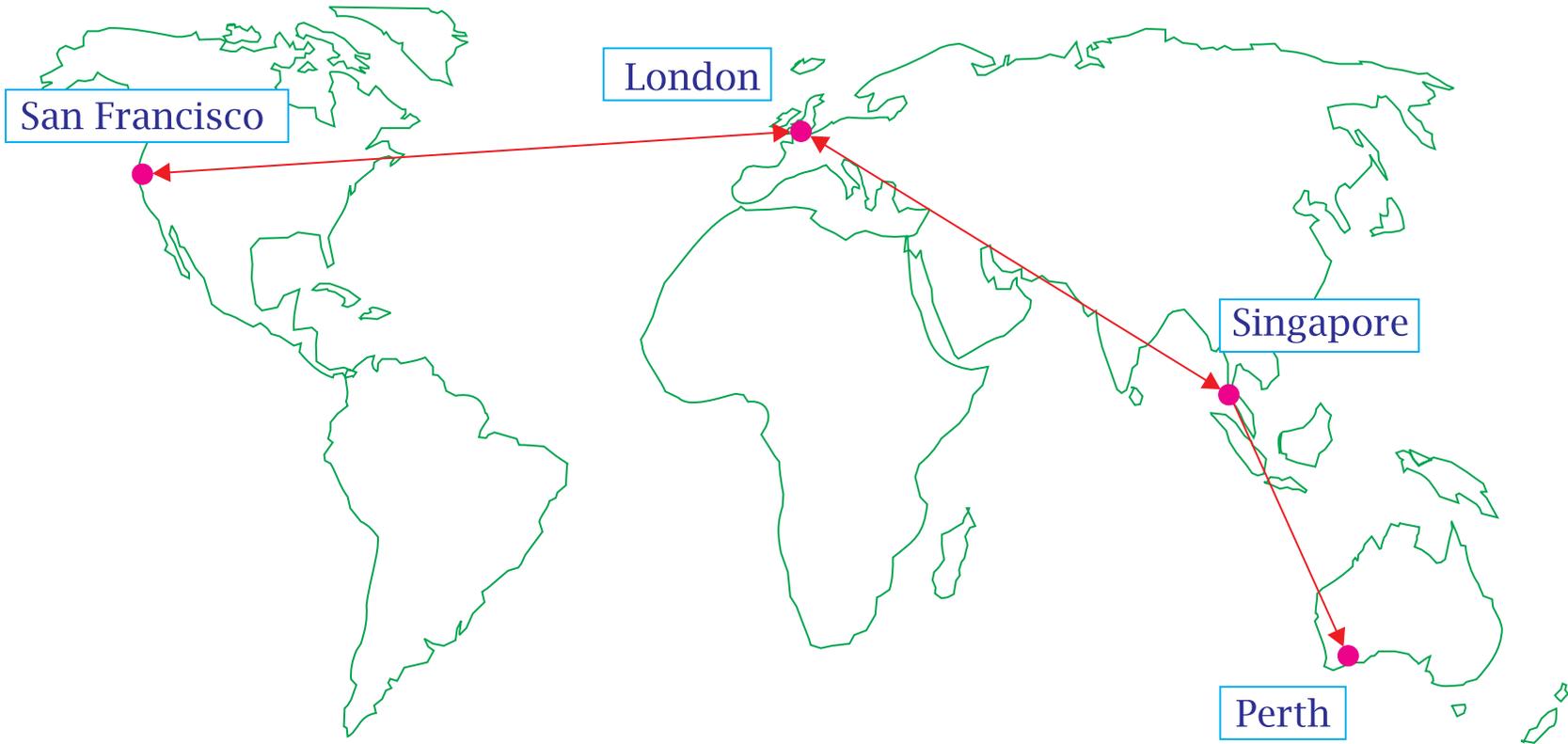
Transitive closures

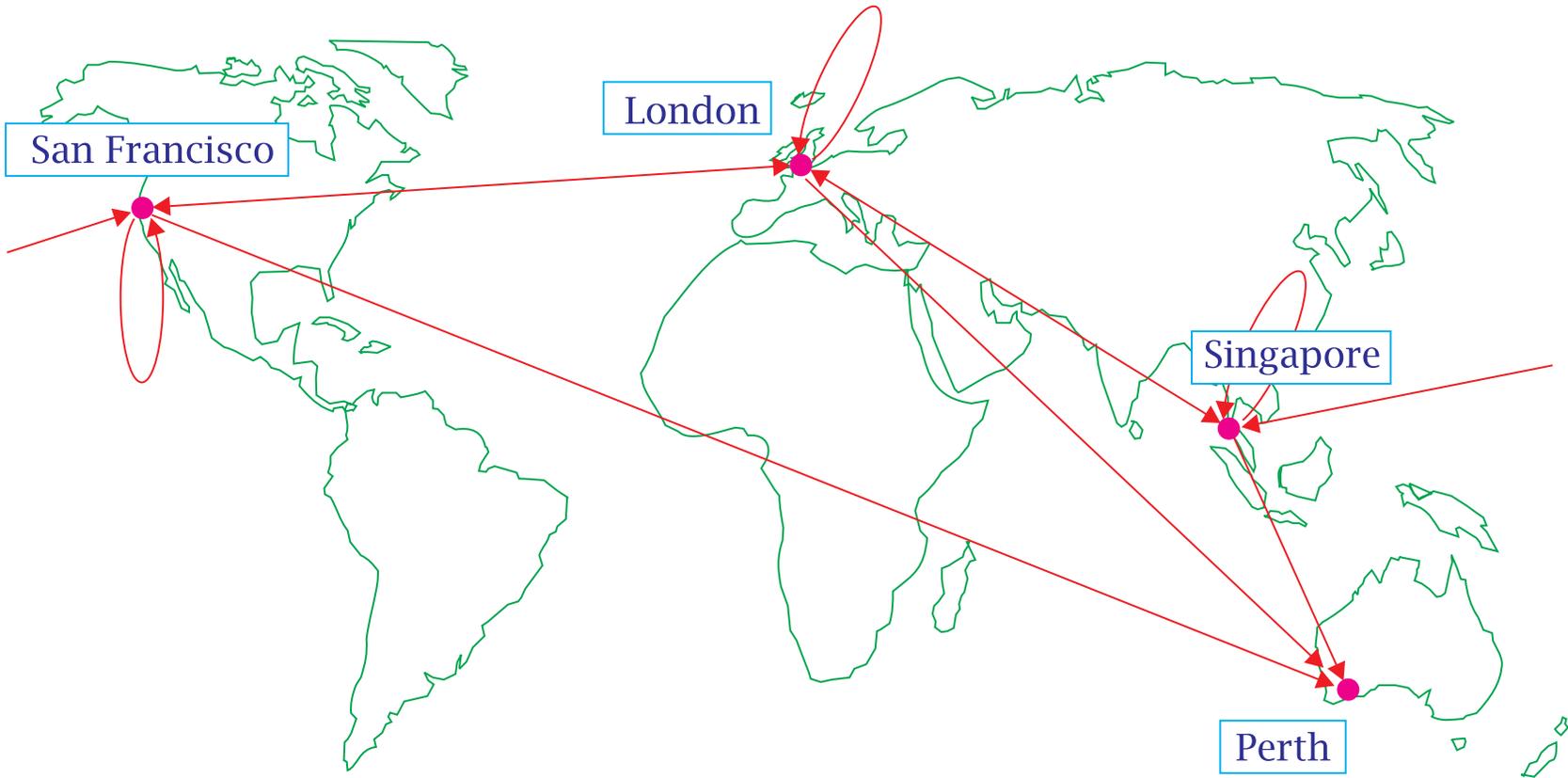
$$R^+ = \bigcup \{ n : \mathbb{N} \mid n \geq 1 \bullet R^n \}$$

$$R^* = R^+ \cup \text{id } X$$

Example

direct = {*singapore* \mapsto *london*, *london* \mapsto *singapore*,
singapore \mapsto *perth*, *london* \mapsto *san_francisco*,
san_francisco \mapsto *london*}





Proof strategy

If \preceq is transitive, then the following rule is valid:

$$\frac{a \preceq b \quad b \preceq c}{a \preceq c} \quad [\preceq \text{ is transitive}]$$

Presentation

We write

$$\begin{array}{c}
 \frac{\overline{a \preceq b} \text{ [reason 1]} \quad \overline{b \preceq c} \text{ [reason 2]}}{\overline{a \preceq c}} \quad \overline{c \preceq d} \text{ [reason 3]} \\
 \hline
 a \preceq d
 \end{array}$$

as

$$\begin{array}{l}
 a \\
 \preceq b \quad \text{[reason 1]} \\
 \preceq c \quad \text{[reason 2]} \\
 \preceq d \quad \text{[reason 3]}
 \end{array}$$

Summary

- relationships between objects
- \leftrightarrow , \mapsto
- dom , ran , \triangleleft , \triangleright , \triangleleft , \triangleright , $(\mid \)$
- \circ , \sim
- reflexivity, symmetry, transitivity
- R^r , S^s , T^+ , T^*
- transitive relations