MANAGED BY UT-BATTELLE FOR THE DEPARTMENT OF ENERGY



Advanced Fault Tolerance Solutions for High Performance Computing

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ORNL East Campus: Site of World Leading Computing and Computational Sciences

Computational Sciences Building

Research Office Building

Engineering Technology Facility

> Old Computational Sciences Building (until June 2003)

Systems Research Team

Joint Institute for Computational Sciences

Research Support Center (Cafeteria, Conference, Visitor)

National Center for Computational Sciences

- 40,000 ft² (3700 m²) computer center:
 - 36-in (~1m) raised floor, 18 ft (5.5 m) deck-to-deck
 - 12 MW of power with 4,800 t of redundant cooling
 - High-ceiling area for visualization lab:
 - 35 MPixel PowerWall, Access Grid, etc.

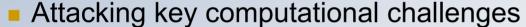


- 2 systems in the Top 500 List of Supercomputer Sites:
 - Jaguar: 10? Cray XT3, MPP with 11500 dual-core Processors ⇒ 119 TFlop.
 - Phoenix: 32? Cray X1E, Vector with 1014 Processors ⇒ 18 TFlop.

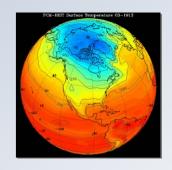


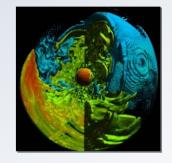
At Forefront in Scientific Computing and Simulation

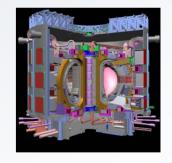
- Leading partnership in developing the National Leadership Computing Facility
 - Leadership-class scientific computing capability
 - □ 100 TFlop/s in 2006/7 (recently installed)
 - □ 250 TFlop/s in 2007/8 (commitment made)
 - 1 PFlop/s in 2008/9 (proposed)



- Climate change
- Nuclear astrophysics
- Fusion energy
- Materials sciences
- Biology
- Providing access to computational resources through high-speed networking (10Gbps)









Talk Outline

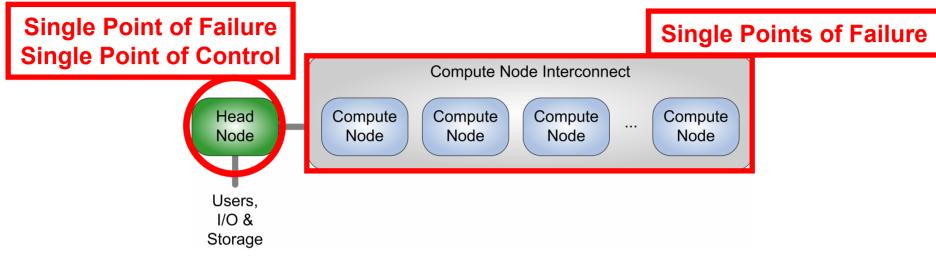
- High performance computing system architectures
- Fault tolerance solutions for head & service nodes:
 - Active/standby with shared storage
 - Active/standby replication
 - Asymmetric active/active replication
 - Symmetric active/active replication
- Fault tolerance solutions for compute nodes:
 - Reactive: Checkpoint/restart and message logging
 - Proactive: Preemptive migration
 - Algorithmic approaches

Advanced Fault Tolerance Solutions for High Performance Computing



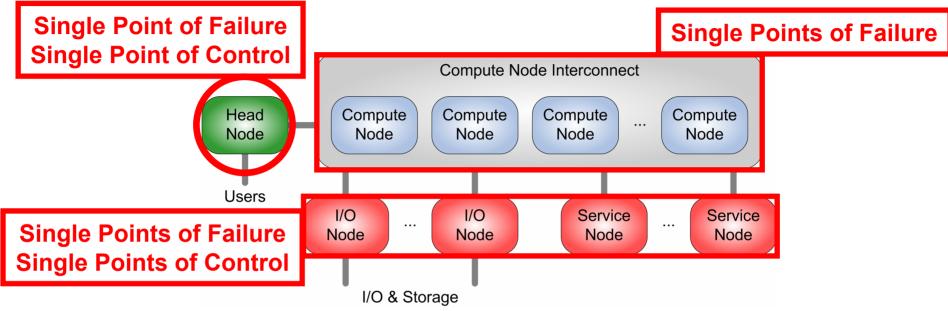
HPC System Architectures

Beowulf Cluster Computing Architecture

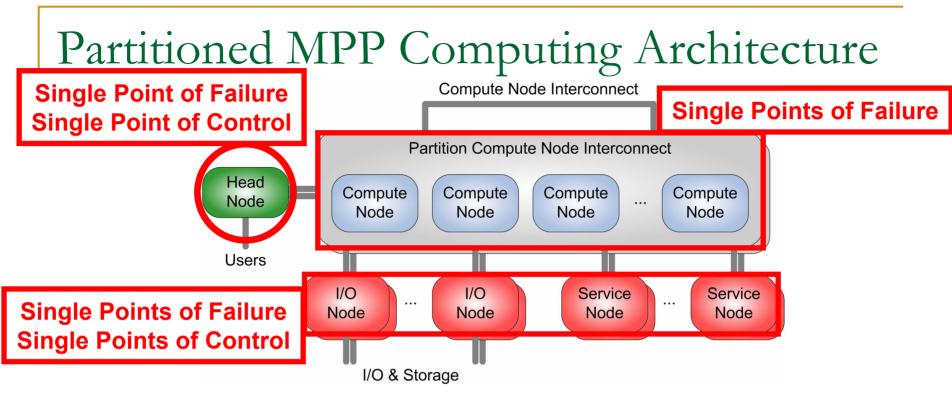


- Single head node manages entire HPC system
- System-wide services are provided by head node:
 - Job & resource management, networked file system, ...
- Local services are provided by compute nodes
 - Message passing (MPI, PVM), ...

Massively Parallel Computing Architecture



- Single head node and additional service nodes manage the entire HPC system
- System-wide services are provided by head node and are offloaded to service nodes, e.g., networked file system
- Local services are provided by service nodes and compute nodes, e.g., message passing



- Single head node manages entire HPC system
- Service nodes manage and support compute nodes belonging to their partitions

Typical Failure Causes in HPC Systems

- Overheating !!!
- Memory and network errors (bit flips)
- Hardware failures due to wear/age of:
 - Hard drives, memory modules, network cards, processors
- Software failures due to bugs in:
 - Operating system, middleware, applications
- → Different scale requires different solutions:
 - → Compute nodes (up to 150,000)
 - → Front-end, service, and I/O nodes (1 to 150)

Availability Measured by the Nines

http://info.nccs.gov/resources - HPC system status at Oak Ridge National Laboratory

9's	Availability	Downtime/Year	Examples
1	90.0%	36 days, 12 hours	Personal Computers
2	99.0%	87 hours, 36 min	Entry Level Business
3	99.9%	8 hours, 45.6 min	ISPs, Mainstream Business
4	99.99%	52 min, 33.6 sec	Data Centers
5	99.999%	5 min, 15.4 sec	Banking, Medical
6	99.9999%	31.5 seconds	Military Defense

- Enterprise-class hardware + Stable Linux kernel = 5+
- Substandard hardware + Good high availability package = 2-3
- Today's supercomputers
 = 1-2
- My desktop = 1-2

Fault Tolerance & High Availability Goals

- Provide high-level Reliability, Availability, and Serviceability (RAS) capabilities
- Eliminate many of the numerous single-points of failure and control in HPC systems

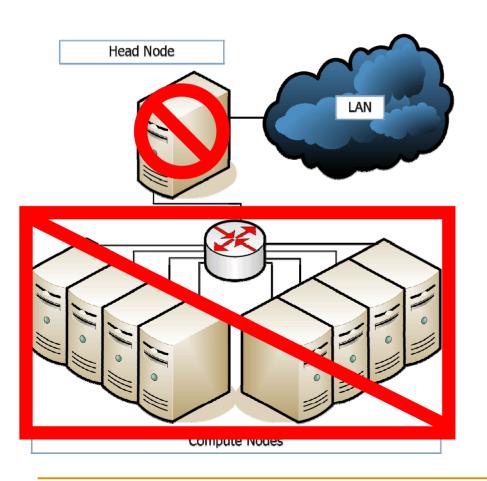
- Development of techniques to enable HPC systems to run computational jobs 24x7 without interruption
- Development of proof-of-concept implementations as blueprint for production-type RAS solutions

Advanced Fault Tolerance Solutions for High Performance Computing



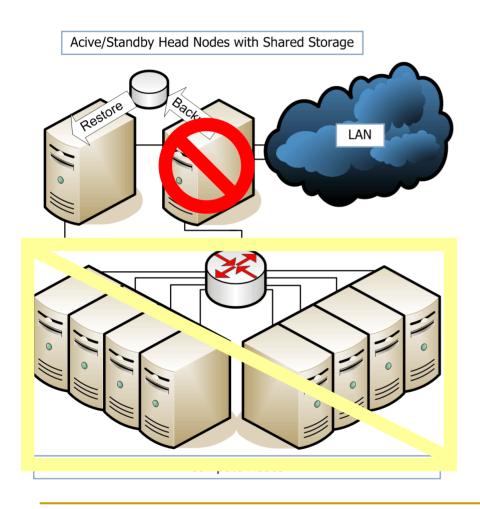
Head and Service Nodes

Single Head/Service Node Problem



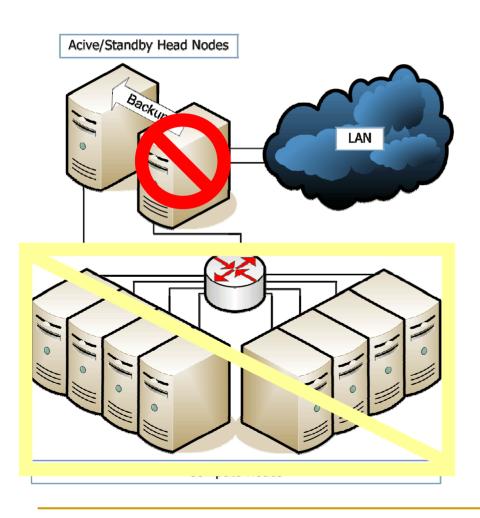
- Single point of failure
- Compute nodes sit idle while head node is down
- A = MTTF / (MTTF + MTTR)
- MTTF depends on head node hardware/software quality
- MTTR depends on the time it takes to repair/replace node
- MTTR = 0 → A = 1.00 (100%)
 continuous availability

Active/Standby with Shared Storage



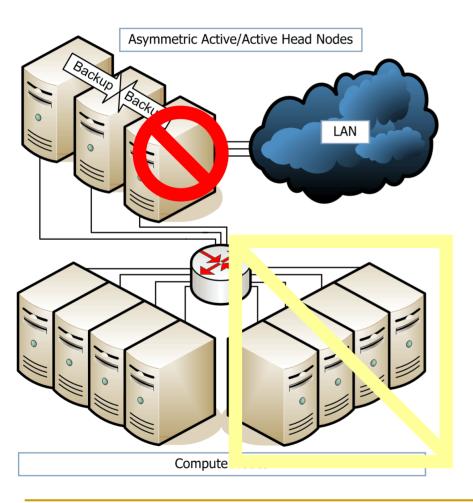
- Single active head node
- Backup to shared storage
- Simple checkpoint/restart
- Fail-over to standby node
- Possible corruption of backup state when failing during backup
- Introduction of a new single point of failure
- Correctness and availability are NOT ALWAYS guaranteed
- SLURM, metadata servers of PVFS and Lustre

Active/Standby Redundancy



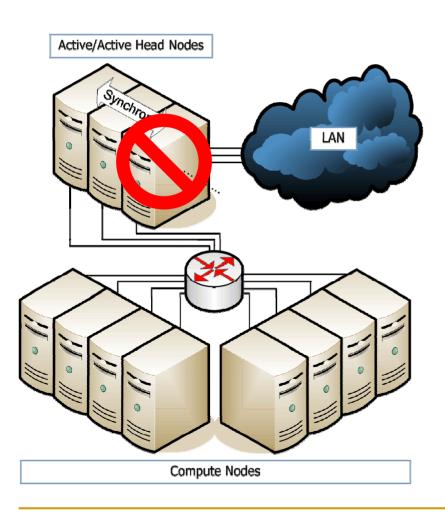
- Single active head node
- Backup to standby node
- Simple checkpoint/restart
- Fail-over to standby node
- Idle standby head node
- Rollback to backup
- Service interruption for failover and restore-over
- Torque on Cray XT
- → HA-OSCAR prototype

Asymmetric Active/Active Redundancy



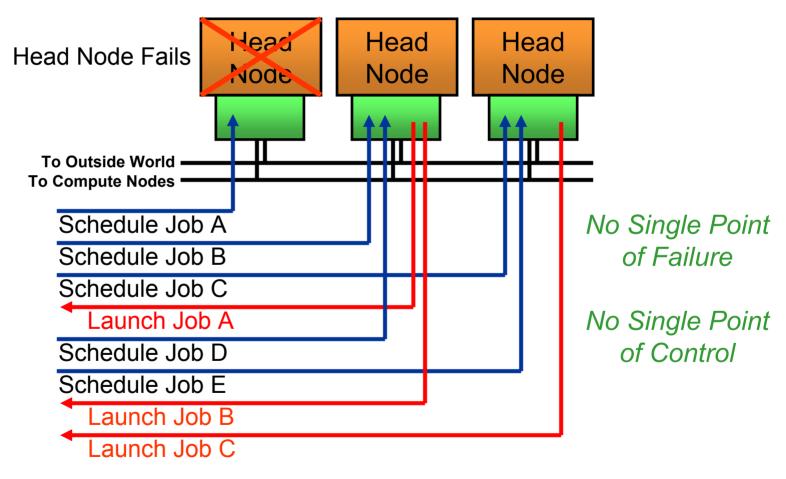
- Many active head nodes
- Work load distribution
- Optional fail-over to standby head node(s) (n+1 or n+m)
- No coordination between active head nodes
- Service interruption for fail-over and restore-over
- Loss of state w/o standby
- Limited use cases, such as high-throughput computing
- Prototype based on HA-OSCAR

Symmetric Active/Active Redundancy

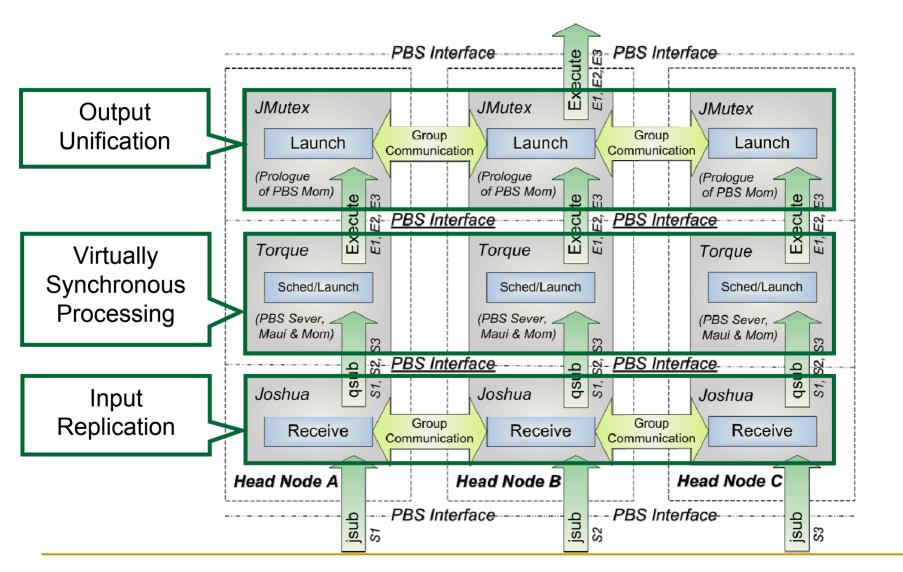


- Many active head nodes
- Work load distribution
- Symmetric replication between head nodes
- Continuous service
- Always up-to-date
- No fail-over, no restore-over
- Virtual synchrony model
- Complex algorithms
- JOSHUA prototype for Torque
- PVFS metadata server

JOSHUA: Symmetric Active/Active Replication for PBS Torque



Symmetric Active/Active Replication



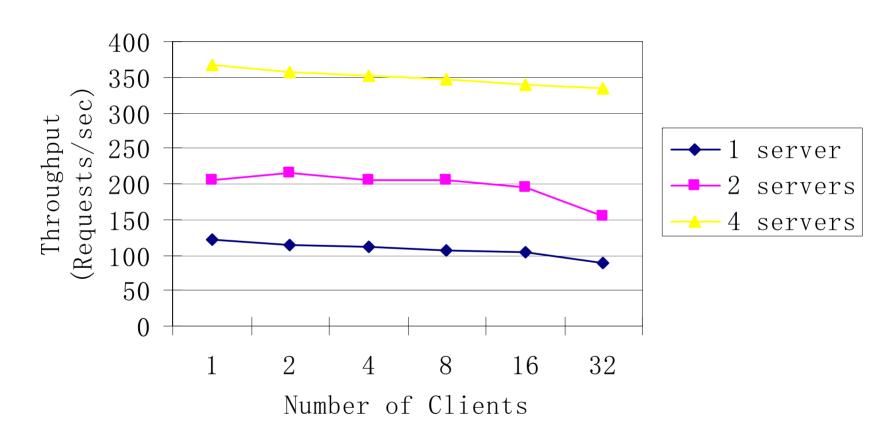
PVFS MDS Performance and Overhead





PVFS MDS Performance and Overhead





Symmetric Active/Active Availability

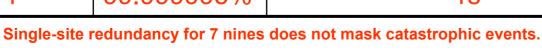
 $A_{\text{system}} = 1 - (1 - A_{\text{component}}) n$

 $T_{down} = 8760 \text{ hours * } (1 - A)$

Single node MTTF: 5000 hours

Single node MTTR: 72 hours

Nodes	Availability	Est. Annual Downtime
1	98.58%	5d 4h 21m
2	99.97%	1h 45m
3	99.9997%	1m 30s
4	99.999995%	1s





Advanced Fault Tolerance Solutions for High Performance Computing



Compute Nodes

Reactive vs. Proactive Fault Tolerance

Reactive fault tolerance:

- State saving during failure-free operation
- State recovery after failure
- Assured quality of service, but limited scalability

Proactive fault tolerance:

- System health monitoring and online reliability modeling
- Failure anticipation and prevention through prediction and reconfiguration before failure
- Highly scalable, but not all failures can be anticipated
- Ideal solution: Matching combination of both

Reactive Fault Tolerance Techniques (1/2)

Checkpoint/restart:

- Application state from all processors is saved regularly on stable storage, such as local disk or networked file system
- On failure, application is restarted using saved state
- Checkpoint always involves data movement (local/network)
- Restart always involves a rollback, i.e., lost computation
- Example: Berkeley Lab Checkpoint/Restart (Linux mod.)
- May be used in combination with message logging to avoid rollback (see next slide)

Reactive Fault Tolerance Techniques (2/2)

Message logging:

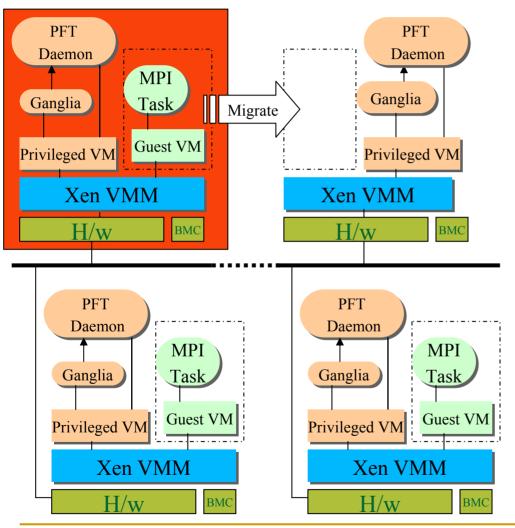
- All messages sent between application processes are logged to a central server
- On failure, only the failed application part is restarted and replayed with saved messages
- Doubles the number of messages
- Message replay involves no rollback
- Example: MPICH-VCL (MPI-based Chandy/Lamport alg.)
- Combination with checkpoint/restart:
 - No rollback / shorter replay time, even higher overhead

Proactive Fault Tolerance Techniques

Preemptive migration:

- System health status is constantly monitored and evaluated
- Monitoring data is processed by a filtering mechanism and/or an online reliability analysis
- Pre-failure indicators are used to predict failures based on current system health status and historic information
- Application parts (processes or virtual machines) are migrated away from compute nodes that are about to fail
- Migration may be performed by stopping the application or live, while keeping the application running

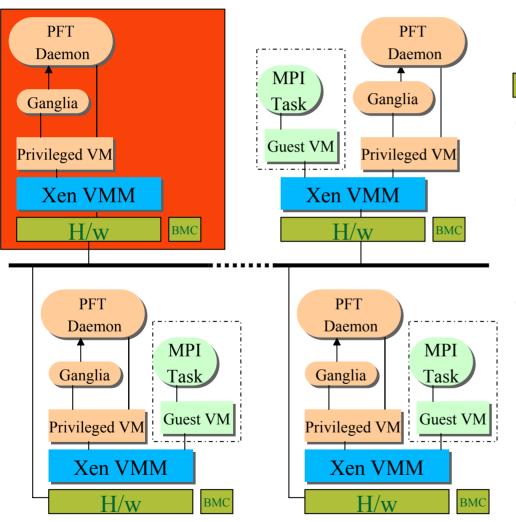
Preemptive Migration with Xen



Baseboard Management Contoller

- Stand-by Xen host, no guest (spare node)
- Deteriorating health >
 migrate guest (w/ MPI app)
 to spare node

Preemptive Migration with Xen



Baseboard Management Contoller

- Stand-by Xen host, no guest (spare node)
- Deteriorating health >
 migrate guest (w/ MPI app)
 to spare node
- Destination host generates unsolicited ARP reply
 - indicates Guest VM IP has moved to new location
 - ARP tells peers to resend packets to new host

Algorithmic Fault Tolerance Approaches

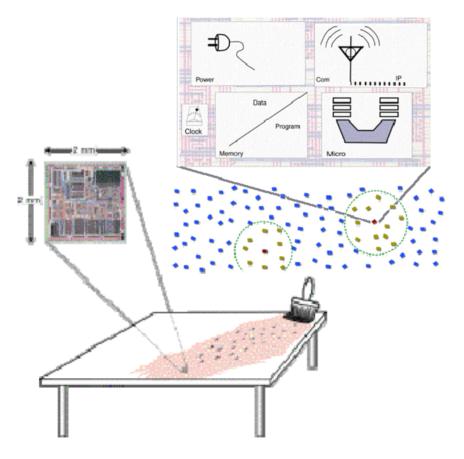
Naturally fault tolerant algorithms

- Processes have only limited knowledge mostly about other processes in their neighborhood
- Application is composed of local algorithms, where a failure has only a minor local impact
- Examples: Chaotic relaxation, peer-to-peer communication

Recovery & erasure codes

- Reconstruction of lost information through algorithmic redundancy within the application
- Rollback to consistent state through reverse computation

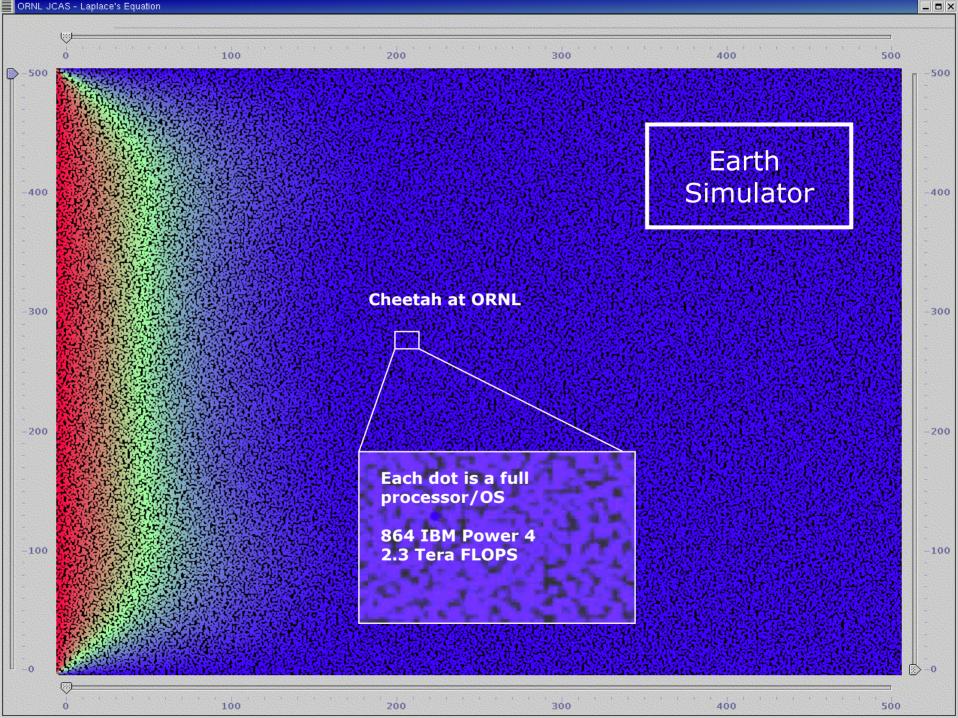
MIT Research: Paintable Computing



- In the future, embedded computers with a radio device will get as small as a paint pigment
- Supercomputers can be easily assembled by just painting a wall of embedded computers
- Applications are driven by cellular algorithms

Cellular Architecture (Smart Dust) Simulator

- Developed at ORNL in Java with native C and Fortran application support using JNI
- Runs as standalone or distributed application
- Lightweight framework simulates up to 1,000,000 lightweight processes on 9 real processors
- Standard and experimental networks:
 - Multi-dimensional mesh/torus
 - Nearest/Random neighbors
- Message driven simulation is not in real-time
- Primitive fault-tolerant MPI support



Summary and Conclusion

- Presented several traditional and advanced fault tolerance technologies for HPC
- Different scale requires different solutions:
 - Compute nodes
 - Front-end, service, and I/O nodes
- Scalable fault tolerance technologies are paramount to the success of large-scale HPC systems

MOLAR: Adaptive Runtime Support for High-end Computing Operating and Runtime Systems

- Addresses the challenges for operating and runtime systems to run large applications efficiently on future ultra-scale high-end computers.
- Part of the <u>Forum to Address Scalable Technology for Runtime</u> and <u>Operating Systems (FAST-OS)</u>.
- MOLAR is a collaborative research effort (<u>www.fastos.org/molar</u>):

















THE SUPERCOMPUTER COMPANY



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