OpenSHMEM Application Programming Interface

http://www.openshmem.org

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Developed by

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Manser, The OpenSHMEM specification belongs to Open Source Software Solutions, Inc. (OSSS), a non-profit organization, under an agreement with SGI. The development work of the specification is supported by the Oak Ridge National Laboratory Extreme Scale Systems Center and the Department of Defense.

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Contents

Introduction

1 The OpenSHMEM Effort

OpenSHMEM is a *Partitioned Global Address Space* (PGAS) library interface specification. OpenSHMEM aims to provide a standard *Application Programming Interface* (API) for SHMEM libraries to aid portability and facilitate uniform predictable results of OpenSHMEM applications by explicitly stating the behavior and semantics of the OpenSHMEM library calls. Through the different versions, OpenSHMEM will continue to address the requirements of the PGAS community. As of this specification, existing vendors are moving towards OpenSHMEM compliant implementations and new vendors are developing OpenSHMEM library implementations to help the users write portable OpenSHMEM code. This ensures that applications can run on multiple platforms without having to deal with subtle vendor-specific implementation differences. For more details on the history of OpenSHMEM please refer to [The](#page-66-0) [History of OpenSHMEM](#page-66-0) section.

The OpenSHMEM¹ effort is driven by the Extreme Scale Systems Center (ESSC) at ORNL and the University of Houston with significant input from the OpenSHMEM community. Besides the specification, the effort also includes providing a reference OpenSHMEM implementation, validation and verification suites, tools, a mailing list and website infrastructure to support specification activities. For more information please refer to: [http://www.openshmem.](http://www.openshmem.org/) [org/](http://www.openshmem.org/).

2 Programming Model Overview

new coase. This entailes mail proplementation can run on manney parameterize, specific implementation differences. For more details on the history of OpenSHMEM pelace refer to TopenSHMEM because with a Chero prossume of t OpenSHMEM implements PGAS by defining remotely accessible data objects as mechanisms to share information among OpenSHMEM processes or *Processing Elements* (PEs) and private data objects that are accessible by the PE itself. The API allows communication and synchronization operations on both private (local) and remotely accessible data objects. The key feature of OpenSHMEM is that data transfer functions are *one-sided* in nature. This means that a local PE executing a data transfer does not require the participation of the remote PE to complete the operation. This allows for overlap between communication and computation to hide data transfer latencies, which makes OpenSHMEM ideal for unstructured, small/medium size data communication patterns. The OpenSHMEM library functions have the potential to provide low-latency, high-bandwidth communication API for use in highly parallelized scalable programs.

The OpenSHMEM interfaces can be used to implement *Single Program Multiple Data* (SPMD) style programs. It provides interfaces to start the OpenSHMEM PEs in parallel, and communication and synchronization interfaces to access remotely accessible data objects across PEs. These interfaces can be leveraged to divide a problem into multiple sub-problems that can solved independently or with coordination using the communication and synchronization interfaces. The OpenSHMEM specification defines library calls, constants, variables, and language bindings for *C* and *Fortran*. The *C++* interface is currently the same as that for *C*. Unlike UPC, Fortran 2008, Titanium, X10 and Chapel, which are all PGAS languages, OpenSHMEM relies on the programmer to use the library calls to implement the correct semantics of its programming model.

An overview of the OpenSHMEM operations is described below:

1. Library Setup and Query

- (a) *Initialization*: The OpenSHMEM library environment is initialized.
- (b) *Query*: The local PE may get number of PEs running the same application and its unique integer identifier.
- (c) *Accessibility*: The local PE can find out if a remote PE is executing the same binary, or if a particular symmetric data object can be accessed by a remote PE, or may obtain a pointer to a symmetric data object on the specified remote PE on shared memory systems.

2. Symmetric Data Object Management

(a) *Allocation*: All executing PEs must participate in the allocation of a symmetric data object with identical arguments.

¹The OpenSHMEM specification is owned by Open Source Software Solutions Inc., a non-profit organization, under an agreement with SGI.

Figure 1: *OpenSHMEM* Memory Model

3 Memory Model

An OpenSHMEM program consists of data objects that are private to each PE and data objects that are remotely accessible by all PEs. Private data objects are stored in the local memory of each PE and can only be accessed by the PE itself; these data objects cannot be accessed by other PEs via OpenSHMEM routines. Private data objects follow the memory model of *C* or *Fortran*. Remotely accessible objects, however, can be accessed by remote PEs using OpenSHMEM routines. Remotely accessible data objects are called *Symmetric Objects*. All symmetric data objects have a corresponding object with the same name, type, size, and offset (from an arbitrary memory address) on all PEs. *Symmetric objects* are accessible by all executing PEs via the OpenSHMEM API. Symmetric data objects accessed via typed OpenSHMEM interfaces are required to be natural aligned based on their type requirements and underlying architecture. In OpenSHMEM the following kinds of data objects are symmetric:

- *Fortran* data objects in common blocks or with the SAVE attribute. These data objects must not be defined in a dynamic shared object (DSO).
- Global and static *C* and *C++* variables. These data objects must not be defined in a DSO.
- *Fortran* arrays allocated with *shpalloc*
- *C* and *C++* data allocated by *shmalloc*

OpenSHMEM dynamic memory allocation routines (*shpalloc* and *shmalloc*) allow collective allocation of *Symmetric Data Objects* on a special memory region called the Symmetric Heap. The Symmetric Heap is created during the execution of a program at a memory location determined by the implementation. The Symmetric Heap may reside on different memory regions on different PEs. Figure [1](#page-6-1) shows how OpenSHMEM implements a PGAS model using remotely accessible (*Symmetric objects*) and private data objects when executing an OpenSHMEM program. Symmetric data objects are stored on the symmetric heap or in the global/static memory section of each PE.

1

4 Execution Model

An OpenSHMEM program consists of a set of OpenSHMEM processes called PEs that execute in a SPMD-like model where each PE can take a different execution path. A PE can be implemented using an OS process or an OS thread^{[2](#page-7-5)}. The PEs progress asynchronously, and can communicate/synchronize via the OpenSHMEM interfaces. All PEs in an OpenSHMEM program should start by calling the initialization function *start*_*pes* before using any of the other OpenSHMEM library routines. As of now, an OpenSHMEM program finishes execution by returning from the main function. On program exit, OpenSHMEM can release all the resources associated to the library.

The PEs of the OpenSHMEM program are identified by unique integers. The identifiers are integers assigned in a monotonically increasing manner from zero to the total number of PEs minus 1. PE identifiers are used for OpenSHMEM calls (e.g. to specify *Put* or *Get* operations on symmetric data objects, collective synchronization calls, etc) or to dictate a control flow for PEs using constructs of *C* or *Fortran*. The identifiers are fixed for the life cycle of the OpenSHMEM program.

4.1 Progress of OpenSHMEM operations

The OpenSHMEM model assumes that computation and communication are naturally overlapped, and that all data transfers eventually complete.

Note to implementors: while delivery can be deferred, for example until a synchronization point at which data is known to be available, high-quality implementations should attempt asynchronous delivery, whenever possible, for performance reasons. Progress will often be ensured through the use of a dedicated progress thread in software, or through network hardware that offloads communication handling from processors, for example.

4.2 Atomicity Guarantees

OpenSHMEM contains a number of routines that operate on symmetric data atomically (Section 8.4). These routines guarantee that accesses by OpenSHMEM's atomic operations will be exclusive, but do not guarantee exclusivity in combination with other routines, either inside OpenSHMEM or outside.

For example: during the execution of a atomic remote integer increment operation on a symmetric variable *X*, no other OpenSHMEM atomic operation may access *X*. After the increment, *X* will have increased its value by *1* on the *target* PE, at which point other atomic operations may then modify that *X*. However, access to the symmetric object *X* with non-atomic operations, such as one-sided *Put* or *Get* operations, will *invalidate* the atomicity guarantees.

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47 48

5 Language Bindings and Conformance

dictate a control flow for PEs using constructs of C or Fortram. The identifiers are fixed for the life cycle

HMEM program.

EHMEM program.

SHMEM operations

SHMEM operations

SHMEM operations

SHMEM operations

start c OpenSHMEM provides ISO *C* and *Fortran* 90 language bindings. Any implementation that provides both *C* and *Fortran* bindings can claim conformance to the specification. An implementation that provides e.g. only a *C* interface may claim to conform to the OpenSHMEM specification with respect to the *C* language, but not to *Fortran*, and should make this clear in its documentation. The OpenSHMEM header files for *C* and *Fortran* must contain only the interfaces and constant names defined in this specification.

OpenSHMEM APIs can be implemented as either functions or macros. However, implementing the interfaces using macros is strongly discouraged as this could severely limit the use of external profiling tools and high-level compiler optimizations. An OpenSHMEM program should avoid defining function names, variables, or identifiers with the prefix *SHMEM*_ (for *C* and *Fortran*), _*SHMEM*_ (for *C*) or with OpenSHMEM API names.

6 Library Constants

Constants Related To Collective Operations

Below are the library constants for collective operations. The constants that start with SHMEM_* are for *Fortran* and _SHMEM_* for *C*. 45 46

²However, implementing a PE using an OS thread requires compiler techniques to implement the OpenSHMEM memory model.

7 Environment Variables

The OpenSHMEM specification provides a set of environment variables that allows users to configure the OpenSHMEM implementation, and receive information about the implementation. The implementations of the specification are free to define additional variables. Currently, the specification defines four environment variables.

8.1.2 SHMEM**_**MY**_**PE

Returns the number of the calling PE.

SYNOPSIS

 C/C_{++} : **int** shmem_my_pe(**void**); **int** _my_pe (**void**);

FORTRAN:

INTEGER SHMEM_MY_PE, ME ME = SHMEM_MY_PE() $ME = MY_PE$ ()

DESCRIPTION

Arguments

None

API description

ances such that
 $\alpha_{\text{V}} = \frac{1}{2}$ ($\alpha_{\text{V}} = \frac{1}{2}$ (and $\alpha_{\text{V}} = \frac{1}{2}$ (b)
 CRAFTION
 C This function returns the PE number of the calling PE. It accepts no arguments. The result is an integer between *0* and *npes* - *1*, where *npes* is the total number of PEs executing the current program.

Return Values

Integer - Between *0* and *npes* - *1*

Notes

Each PE has a unique number or identifier.

EXAMPLES

The following _*my*_*pe* example is for *C/C++* programs:

#include <stdio.h> **#include** <shmem.h>

int main(**void**) {

int me;

```
start_pes(0);
me = my_pe();printf("My PE id is: %d\n", me);
return 0;
```
8.1.3 SHMEM**_**N**_**PES

Returns the number of PEs running in an application.

SYNOPSIS

}

C/C++: **int** shmem_n_pes(**void**); **int** _num_pes (**void**);

FORTRAN:

Arguments

IN *pe* Specific pe that needs to be checked if accessible from the local PE.

API description

accouss mutup particulare through an initial concerned noise and the present mutual concerned noise was applications using OpenSHMEM on an Altix, full OpenSHMEM support is whallable between processes executable files is s *shmem*_*pe*_*accessible* is a query function that indicates whether a specified PE is accessible via OpenSHMEM from the local PE. The *shmem*_*pe*_*accessible* function returns *TRUE* only if the remote PE is a process running from the same executable file as the local PE, indicating that full OpenSHMEM support for symmetric data objects (that resides in the static memory and symmetric heap) is available, otherwise it returns *FALSE*. This function may be particular useful for hybrid programming with other communication libraries (such as a MPI) or parallel languages. For example, on SGI Altix series systems, OpenSHMEM is supported across multiple partitioned hosts and InfiniBand connected hosts. When running multiple executable MPI applications using OpenSHMEM on an Altix, full OpenSHMEM support is available between processes running from the same executable file. However, OpenSHMEM support between processes of different executable files is supported only for data objects on the symmetric heap, since static data objects are not symmetric between different executable files.

Return Values

C: The return value is 1 if the specified PE is a valid remote PE for OpenSHMEM functions; otherwise,it is 0.

Fortran: The return value is *.TRUE*. if the specified PE is a valid remote PE for OpenSHMEM functions; otherwise, it is *.FALSE.*.

Notes

None.

8.1.5 SHMEM**_**ADDR**_**ACCESSIBLE

Determines whether an address is accessible via OpenSHMEM data transfers operations from the specified remote PE.

SYNOPSIS

```
C/C++
```
int shmem_addr_accessible(**void** *addr, **int** pe);

FORTRAN:

LOGICAL LOG, SHMEM_ADDR_ACCESSIBLE **INTEGER** pe LOG = SHMEM_ADDR_ACCESSIBLE(addr, pe)

DESCRIPTION

API description

*shmem*_*addr*_*accessible* is a query function that indicates whether a local address is accessible via OpenSHMEM 44 operations from the specified remote PE. 45

This function verifies that the data object is symmetric and accessible with respect to a remote PE via OpenSHMEM data transfer functions. The specified address *addr* is a data object on the local PE.

TThis function may be particular useful for hybrid programming with other communication libraries (such as a MPI) or parallel languages. For example, in SGI Altix series systems, for multiple executable MPI

1


```
\frac{1}{2} and \frac{1}{2} a
PROGRAM REMOTEWRITE
INCLUDE 'shmem.fh'
INTEGER BIGD(100)
SAVE BIGD
INTEGER POINTEE(*)
POINTER (PTR,POINTEE)
CALL START_PES(0)
IF (MY_PE() .EQ. 0) THEN
   ! initialize PE 1's BIGD array
   PTR = SHMEM_PTR(BIGD, 1) ! get address of PE 1's BIGD
                                        ! array
   DO I=1,100
         POINTEE(I) = IENDDO
ENDIF
CALL SHMEM_BARRIER_ALL
IF (MY_PE() .EQ. 1) THEN
   PRINT*,'BIGD on PE 1 is: '
   PRINT*,BIGD
ENDIF
END
This is the equivalent program written in C:
#include <shmem.h>
int main(void)
{
   static int bigd[100];
       int *ptr;
       int i;
   start_pes(0);
   if (\text{my_pe}) == 0 {
       /* initialize PE 1's bigd array */
       ptr = shmem_ptr(bigd, 1);
       for (i=0; i<100; i++)
           *ptr++ = i+1;}
    shmem_barrier_all();
   if (\text{my_pe}) == 1)printf("bigd on PE 1 is:\n");
       for (i=0; i<100; i++)
          printf("d\nmid n", bigd[i]);
       printf(\sqrt[m]{n^n});
    }
  return 1;
}
```
8.2 Memory Management Operations

OpenSHMEM provides a set of APIs for managing the symmetric heap. The APIs allow to dynamically allocate, deallocate, reallocate and align symmetric data objects in the symmetric heap, in *C* and *Fortran*.

8.2.1 SHMALLOC, SHFREE, SHREALLOC, SHMEMALIGN

Symmetric heap memory management functions.

SYNOPSIS

DESCRIPTION

Arguments

API description

The *shmalloc* function returns a pointer to a block of at least size bytes suitably aligned for any use. This space is allocated from the symmetric heap (in contrast to *malloc*, which allocates from the private heap). The *shmemalign* function allocates a block in the symmetric heap that has a byte alignment specified by the alignment argument.

- The *shfree* function causes the block to which *ptr* points to be deallocated, that is, made available for further allocation. If ptr is a null pointer, no action occurs; otherwise, if the argument does not match a pointer earlier returned by a symmetric heap function, or if the space has already been deallocated, *malloc*_*error* is set to indicate the error, and *shfree* returns. α 21 22
- IN size In bytes, to request a block to be allocated from the symmetric hero.

This argument is of type size, the control control particle heap.

This argument is of type size, the control of the symmetric heap.

IN prime The *shrealloc* function changes the size of the block to which ptr points to the size (in bytes) specified by size. The contents of the block are unchanged up to the lesser of the new and old sizes. If the new size is larger, the value of the newly allocated portion of the block is indeterminate. If *ptr* is a *NULL* pointer, the *shrealloc* function behaves like the *shmalloc* function for the specified size. If size is *0* and ptr is not a *NULL* pointer, the block to which it points is freed. Otherwise, if ptr does not match a pointer earlier returned by a symmetric heap function, or if the space has already been deallocated, the *malloc*_*error* variable is set to indicate the error, and *shrealloc* returns a *NULL* pointer. If the space cannot be allocated, the block to which ptr points is unchanged. 23 24 25 26 27 28 29 30
	- The *shmalloc*, *shfree*, and *shrealloc* functions are provided so that multiple PEs in an application can allocate symmetric, remotely accessible memory blocks. These memory blocks can then be used with OpenSHMEM communication routines. Each of these functions call the *shmem*_*barrier*_*all* function before returning; this ensures that all PEs participate in the memory allocation, and that the memory on other PEs can be used as soon as the local PE returns. The user is responsible for calling these functions with identical argument(s) on all PEs; if differing size arguments are used, subsequent calls may not return the same symmetric heap address on all PEs.

Return Values

- The *shmalloc* function returns a pointer to the allocated space (which should be identical on all PEs); otherwise, it returns a *NULL* pointer (with *malloc*_*error* set).
	- The *shfree* function returns no value.
		- The *shrealloc* function returns a pointer to the allocated space (which may have moved); otherwise, it returns a null pointer (with *malloc*_*error* set).

Notes

hmalloc, shfree, and *shrealloc* functions differ from the private heap allocation functions in that all PEs in an application must call them (a barrier is used to ensure this). 47 48

8.2.2 SHPALLOC

Allocates a block of memory from the symmetric heap.

SYNOPSIS

```
FORTRAN:
POINTER (addr, A(1))
INTEGER (length, errcode, abort)
CALL SHPALLOC(addr, length, errcode, abort)
```
DESCRIPTION

Arguments

API description

SHPALLOC allocates a block of memory from the program's symmetric heap that is greater than or equal to the size requested. To maintain symmetric heap consistency, all PEs in an program must call *SHPALLOC* with the same value of length; if any PEs are missing, the program will hang.

By using the *Fortran POINTER* mechanism in the following manner, you can use array *A* to refer to the block allocated by *SHPALLOC*: *POINTER* (*addr*, *A*())

Return Values

Notes

The total size of the symmetric heap is determined at job startup. One may adjust the size of the heap using the *SMA*_*SYMMETRIC*_*SIZE* environment variable (if available).

8.2.3 SHPCLMOVE

Extends a symmetric heap block or copies the contents of the block into a larger block.

SYNOPSIS

```
FORTRAN:
POINTER (addr, A(1))
INTEGER length, status, abort
CALL SHPCLMOVE (addr, length, status, abort)
```
DESCRIPTION

IN *abort* Abort code. Nonzero requests abort on error; *0* requests an error code.

API description

The *SHPCLMOVE* function either extends a symmetric heap block if the block is followed by a large enough free block or copies the contents of the existing block to a larger block and returns a status code indicating that the block was moved. This function also can reduce the size of a block if the new length is less than the old length. All PEs in a program must call *SHPCLMOVE* with the same value of *addr* to maintain symmetric heap consistency; if any PEs are missing, the program hangs.

Return Values

Notes

None.

8.2.4 SHPDEALLC

Returns a memory block to the symmetric heap.

SYNOPSIS

DESCRIPTION

API description

SHPDEALLC returns a block of memory (allocated using *SHPALLOC*) to the list of available space in the symmetric heap. To maintain symmetric heap consistency, all PEs in a program must call *SHPDEALLC* with the same value of *addr*; if any PEs are missing, the program hangs.

Return Values

8. OPENSHMEM LIBRARY API 15

Notes

None.

8.3 Remote Memory Access Operations

Remote Memory Access (RMA) operations described in this section are one-sided communication mechanisms of the OpenSHMEM API. While using these mechanisms, the programmer is required to provide parameters only on the calling side. A characteristic of one-sided communication is that it decouples communication from the synchronization. One-sided communication mechanisms transfer the data but do not synchronize the sender of the data with the receiver of the data.

OpenSHMEM RMA operations are all performed on the symmetric objects. The initiator PE of the call is designated as *source*, and the PE in which memory is accessed is designated as *target*. In the case of the remote update operation, *Put*, the origin is the *source* PE and the destination PE is the *target* PE. In the case of the remote read operation, *Get*, the origin is the *target* PE and the destination is the *source* PE.

OpenSHMEM provides three different types of one-sided communication interfaces. *shmem*_*put*<*bits*> interface transfers data in chunks of bits. *shmem*_*put32*, for example, copies data to a *target* PE in chunks of 32 bits. *shmem*_<*datatype*>_*put* interface copies elements of type *datatype* from a *source* PE to a *target* PE. For example, *shmem*_*integer*_*put*, copies elements of type integer from a *source* PE to a *target* PE. *shmem*_<*datatype*>_*p* interface is similar to *shmem*_<*datatype*>_*put* except that it only transfers one element of type *datatype*.

OpenSHMEM provides interfaces for transferring both contiguous and non-contiguous data. The non-contiguous data transfer interfaces are prefixed with "*i*". *shmem*_<*datatype*>_*iput* interface, for example, copies strided data elements from the *source* PE to a *target* PE.

8.3.1 SHMEM**_**PUT

The put routines provide a method for copying data from a contiguous local data object to a data object on a specified PE.

SYNOPSIS

C/C_{++} :

```
HMEM API. While using these mechanisms, the programer is required to provide parameters only on<br>
\alpha. A characteristic of one-sided communication is that it decouples communication from the synchronization<br>
conmunication
void shmem_double_put(double target, const double *source, size_t len, int pe);
void shmem_float_put(float *target, const float *source, size_t len, int pe);
void shmem_int_put(int *target, const int *source, size_t len, int pe);
void shmem_long_put(long *target, const long *source, size_t len, int pe);
void shmem_longdouble_put(long double *target, const long double *source, size_t len,int pe);
void shmem_longlong_put(long long *target, const long long *source, size_t len, int pe);
void shmem_put32(void *target, const void *source, size_t len, int pe);
void shmem_put64(void *target, const void *source, size_t len, int pe);
void shmem_put128(void *target, const void *source, size_t len, int pe);
void shmem_putmem(void *target, const void *source, size_t len, int pe);
void shmem_short_put(short*target, const short*source, size_t len, int pe);
FORTRAN:
CALL SHMEM_CHARACTER_PUT(target, source, len, pe)
```
CALL SHMEM_COMPLEX_PUT(target, source, len, pe) **CALL** SHMEM_DOUBLE_PUT(target, source, len, pe) **CALL** SHMEM_INTEGER_PUT(target, source, len, pe) **CALL** SHMEM_LOGICAL_PUT(target, source, len, pe) **CALL** SHMEM_PUT(target, source, len, pe) **CALL** SHMEM_PUT4(target, source, len, pe)

DESCRIPTION

API description

The routines return after the data has been copied out of the *source* array on the local PE. The delivery of data words into the data object on the destination PE may occur in any order. Furthermore, two successive put operations may deliver data out of order unless a call to *shmem*_*fence* is introduced between the two calls.

The *target* and *source* data objects must conform to certain typing constraints, which are as follows:

None.

Notes

If you are using *Fortran*, data types must be of default size. For example, a real variable must be declared

as *REAL*, *REAL*4*, or *REAL(KIND=4)*.

EXAMPLES

The following *shmem*_*put* example is for programs:

```
#include <stdio.h>
#include <shmem.h>
int main(void)
{
   long source[10] = { 1, 2, 3, 4, 5,
                       6, 7, 8, 9, 10 };
   static long target [10];
  start_pes(0);
   if (\text{my_pe}) == 0 {
      /* put 10 words into target on PE 1 */shmem_long_put(target, source, 10, 1);
   }
   shmem_barrier_all(); /* sync sender and receiver */printf("target[0] on PE %d is %d\n", _my_pe(), target[0]);
   return 1;
}
```
8.3.2 SHMEM**_**P

Copies one data item to a remote PE.

SYNOPSIS

$C/C++$:

```
start_peach (i) \mu purinted lines and resident the control of the purinted and resident the plane in the particular particular of the particular particular of the plane of the plane of the plane of the control of the co
void shmem_char_p(char *addr, char value, int pe);
void shmem_short_p(short *addr, short value, int pe);
void shmem_int_p(int *addr, int value, int pe);
void shmem_long_p(long *addr, long value, int pe);
void shmem_longlong_p(long long *addr, long long value, int pe);
void shmem_float_p(float *addr, float value, int pe);
void shmem_double_p(double *addr, double value, int pe);
void shmem_longdouble_p(long double *addr, long double value, int pe);
```
DESCRIPTION

API description

These routines provide a very low latency put capability for single elements of most basic types. As with *shmem*_*put*, these functions start the remote transfer and may return before the data is delivered to the remote PE. Use *shmem*_*quiet* to force completion of all remote *Put* transfers.

EXAMPLES

```
when (NBI1927)<br>
except (double ) shmallec(sizeof (+f));<br>
except = _Figure(1)<br>
if \sin a - 0<br>
infinite.comble_p(f, e, i);<br>
if \sin a - 0<br>
infinite.comble_p(f, e, i);<br>
if \sin a - 1<br>
if \sin a -The following simple example uses shmem_double_p in a C program.
      #include <stdio.h>
      #include <math.h>
      #include <shmem.h>
      static const double e = 2.71828182;
      static const double epsilon = 0.00000001;
      int main(void)
      {
         double *f;
         int me;
         start_pes(0);
         me = my_pe();
         f = (double *) shmalloc(sizeof (*f));
         *f = 3.1415927;shmem_barrier_all();
         if (me == 0)
             shmem_double_p(f, e, 1);
         shmem barrier all();
         if (me == 1)
             printf("%s\n", (fabs (*f - e) < epsilon) ? "OK" : "FAIL");
         return 0;
      }
8.3.3 SHMEM_IPUT
Copies strided data to a specified PE.
SYNOPSIS
      C/C++void shmem_double_iput(double *target, const double *source, ptrdiff_t tst, ptrdiff_t sst,
           size_t nelems, int pe);
      void shmem_float_iput(float *target, const float *source, ptrdiff_t tst, ptrdiff_t sst,
           size_t nelems, int pe);
      void shmem_int_iput(int *target, const int *source, ptrdiff_t tst, ptrdiff_t sst, size_t
          nelems, int pe);
      void shmem_iput32(void *target, const void *source, ptrdiff_t tst, ptrdiff_t sst, size_t
           nelems, int pe);
      void shmem_iput64(void *target, const void *source, ptrdiff_t tst, ptrdiff_t sst, size_t
          nelems, int pe);
```
- **void** shmem_iput128(**void** *target, **const void** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe);
	- **void** shmem_long_iput(**long** *target, **const long** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe);
- **void** shmem_longdouble_iput(**long double** *target, **const long double** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe);
- **void** shmem_longlong_iput(**long long** *target, **const long long** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe);
- **void** shmem_short_iput(**short** *target, **const short** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe);

FORTRAN:

- **INTEGER** tst, sst, nelems, pe 47
- **CALL** SHMEM_COMPLEX_IPUT(target, source, tst, sst, nelems, pe) 48

DESCRIPTION

API description

The *iput* routines provide a method for copying strided data elements (specified by *sst*) of an array from a *source* array on the local PE to locations specified by stride *tst* on a *target* array on specified remote PE. Both strides, *tst* and *sst* must be greater than or equal to *1*. The routines return when the data has been copied out of the *source* array on the local PE but not necessarily before the data has been delivered to the remote data object.

The *target* and *source* data objects must conform to typing constraints, which are as follows:


```
CALL SHMEM_DOUBLE_GET(target, source, nelems, pe)
CALL SHMEM_GET4(target, source, nelems, pe)
CALL SHMEM_GET8(target, source, nelems, pe)
CALL SHMEM_GET32(target, source, nelems, pe)
CALL SHMEM_GET128(target, source, nelems, pe)
CALL SHMEM_GETMEM(target, source, nelems, pe)
CALL SHMEM_INTEGER_GET(target, source, nelems, pe)
CALL SHMEM_LOGICAL_GET(target, source, nelems, pe)
CALL SHMEM_REAL_GET(target, source, nelems, pe)
```
DESCRIPTION

API description

The get routines provide a method for copying a contiguous symmetric data object from a different PE to a contiguous data object on a the local PE. The routines return after the data has been delivered to the *target* array on the local PE.

The *target* and *source* data objects must conform to typing constraints, which are as follows:

Return Values

None.


```
int main(void)
{
   int me, npes;
  long y = -1;
  start pes(0);
  me = my_pe();
  npes = \text{num}pes();if (me == 0)
      y = shmem_long_g(&x, 1);
   printf("%d: y = %Id\n\infty", me, y);
   return 0;
}
```
8.3.6 SHMEM**_**IGET

Copies strided data from a specified PE.

SYNOPSIS

DRAFT C/C_{++} : **void** shmem_double_iget(**double** *target, **const double** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); **void** shmem_float_iget(**float** *target, **const float** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); **void** shmem_iget32(**void** *target, **const void** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); **void** shmem_iget64(**void** *target, **const void** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); **void** shmem_iget128(**void** *target, **const void** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); **void** shmem_int_iget(**int** *target, **const int** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); **void** shmem_long_iget(**long** *target, **const long** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); **void** shmem_longdouble_iget(**long double** *target, **const long double** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); **void** shmem_longlong_iget(**long long** *target, **const long long** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); **void** shmem_short_iget(**short** *target, **const short** *source, ptrdiff_t tst, ptrdiff_t sst, **size_t** nelems, **int** pe); FORTRAN: **INTEGER** tst, sst, nelems, pe **CALL** SHMEM_COMPLEX_IGET(target, source, tst, sst, nelems, pe) **CALL** SHMEM_DOUBLE_IGET(target, source, tst, sst, nelems, pe) **CALL** SHMEM_IGET4(target, source, tst, sst, nelems, pe) **CALL** SHMEM_IGET8(target, source, tst, sst, nelems, pe) **CALL** SHMEM_IGET32(target, source, tst, sst, nelems, pe) **CALL** SHMEM_IGET64(target, source, tst, sst, nelems, pe) **CALL** SHMEM_IGET128(target, source, tst, sst, nelems, pe) **CALL** SHMEM_INTEGER_IGET(target, source, tst, sst, nelems, pe) **CALL** SHMEM_LOGICAL_IGET(target, source, tst, sst, nelems, pe)

CALL SHMEM_REAL_IGET(target, source, tst, sst, nelems, pe)

DESCRIPTION


```
PROGRAM STRIDELOGICAL
LOGICAL SOURCE(10), TARGET(5)
SAVE SOURCE ! SAVE MAKES IT REMOTELY ACCESSIBLE
DATA SOURCE /.T.,.F.,.T.,.F.,.T.,.F.,.T.,.F.,.T.,.F./
DATA TARGET / 5*.F. /
CALL START_PES(2)
IF (MY_PE() .EQ. 0) THEN
  CALL SHMEM_LOGICAL_IGET(TARGET, SOURCE, 1, 2, 5, 1)
   PRINT*,'TARGET AFTER SHMEM_LOGICAL_IGET:',TARGET
ENDIF
CALL SHMEM_BARRIER_ALL
```
8.4 Atomic Memory Operations

Atomic Memory Operation (AMO) is a one-sided communication mechanism that combines memory update operations with atomicity guarantees described in Section 4.2. Similar to the RMA routines, described in Section 8.3, the AMOs are performed only on symmetric objects. OpenSHMEM defines the two types of AMO routines:

• The *fetch-and-operate* routines combine memory update and fetch operations in a single atomic operation. The routines return after the data has been fetched and delivered to the local PE.

```
The fetch-and-operate operations include: SHMEM_CSWAP, SHMEM_SWAP, SHMEM_FINC, and
SHMEM_FADD.
```
imic Memory Operations

immory Operations

comorginal (AMO) is a one-sided communication mechanism that combines memory update operation

icity guarantes described in Section 4.2. Similar to the RMA routines, described in • The *non-fetch* atomic routines update the remote memory in a single atomic operation. A *non-fetch* atomic routine starts the atomic operation and may return before the operation execution on the remote PE. To force completion for these *non-fetch* atomic routines, *shmem*_*quiet*, *shmem*_*barrier*, or *shmem*_*barrierall* can be used by an OpenSHMEM program.

The *non-fetch* operations include: *SHMEM*_*INC* and *SHMEM*_*ADD*.

8.4.1 SHMEM**_**ADD

Performs an atomic add operation on a remote symmetric data object.

SYNOPSIS

$C/C++$

```
void shmem_int_add(int *target, int value, int pe);
void shmem_long_add(long *target, long value, int pe);
void shmem_longlong_add(long long *target, long long value, int pe);
```
FORTRAN:

```
INTEGER pe
CALL SHMEM_INT4_ADD(target, value, pe)
CALL SHMEM_INT8_ADD(target, value, pe)
```
DESCRIPTION

IN *pe* An integer that indicates the PE number upon which *target* is to be updated. If you are using *Fortran*, it must be a default integer value.

API description

The *shmem*_*add* routine performs an atomic add operation. It adds value to *target* on PE *pe* and atomically increments the *target* without returning the value.

Return Values

None.

Notes

The term remotely accessible is defined in the Introduction.

EXAMPLES

```
The term remotely accessible is defined in the Introduction.<br>
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main (woid)<br>
int me, cld;<br>
main (woid)<br>
int me, cld;<br>
is x - xy = 0;<br>
claim the main (woid)<br>
is x + y = 0;<br>
claim the main interaction of the subset of the s
int main(void)
{
    int me, old;
    start_pes(0);
    me = my_pe();
    old = -1;dst = 22;shmem_barrier_all();
    if (me == 1) {
         old = shmem\_int\_fadd(@dst, 44, 0);}
    shmem_barrier_all();
    printf("%d: old = %d, dst = %d\n", me, old, dst);
    return 0;
}
```
8.4.2 SHMEM**_**CSWAP

Performs an atomic conditional swap to a remote data object. 30

SYNOPSIS

31 32 33

$C/C++$

```
int shmem_int_cswap(int *target, int cond, int value, int pe);
           long shmem_long_cswap(long *target, long cond, long value, int pe);
           long shmem_longlong_cswap(long long *target, long long cond, long long value, int pe);
           FORTRAN:
           INTEGER pe
           INTEGER(KIND=4) SHMEM_INT4_CSWAP
           ires = SHMEM_INT4_CSWAP(target, cond, value, pe)
           INTEGER(KIND=8) SHMEM_INT8_CSWAP
           ires = SHMEM_INT8_CSWAP(target, cond, value, pe)
     DESCRIPTION
           Arguments
                 OUT target The remotely accessible integer data object to be updated on the remote
                                               PE.
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```


API description

The conditional swap routines conditionally update a *target* data object on an arbitrary PE and return the prior contents of the data object in one atomic operation.

The *target* and source data objects must conform to certain typing constraints, which are as follows:

Return Values

The contents that had been in the *target* data object on the remote PE prior to the conditional swap. Data type is the same as the *target* data type.

Notes

None.

EXAMPLES

The following call ensures that the first PE to execute the conditional swap will successfully write its PE number to *race*_*winner* on PE *0*.

```
int main(void)
{
   static int race_winner = -1;
  int oldval;
  start_pes(2);
  oldval = shmem_int_cswap(&race_winner, -1, _my_pe(), 0);
   if(oldval == -1) printf("pe %d was first\n",_my_pe());
   return 1;
}
```
8.4.3 SHMEM**_**SWAP

Performs an atomic swap to a remote data object.

SYNOPSIS

 $C/C++$: **double** shmem_double_swap(**double** *target, **double** value, **int** pe); **float** shmem_float_swap(**float** *target, **float** value, **int** pe); **int** shmem_int_swap(**int** *target, **int** value, **int** pe); **long** shmem_long_swap(**long** *target, **long** value, **int** pe); **long long** shmem_longlong_swap(**long long** *target, **long long** value, **int** pe); **long** shmem_swap(**long** *target, **long** value, **int** pe);

FORTRAN:


```
start_pes(0);
  me = my_pe();npes = num_pes();target = (long *) shmalloc(sizeof (*target));
  *target = me;
  shmem barrier all();
  new_val = me;
  if (me & 1){
     swapped_val = shmem_long_swap(target, new_val, (me + 1) % npes);
     printf("%d: target = %d, swapped = %d\n", me, *target, swapped_val);
   }
  shfree(target);
  return 0;
}
```
8.4.4 SHMEM**_**FINC

Performs an atomic fetch-and-increment operation on a remote data object.

SYNOPSIS

$C/C++$:

```
int shmem_int_finc(int *target, int pe);
long shmem_long_finc(long *target, int pe);
long long shmem_longlong_finc(long long *target, int pe);
```
FORTRAN:

```
INTEGER pe
INTEGER(KIND=4) SHMEM_INT4_FINC, target4
INTEGER(KIND=8) SHMEM_INT8_FINC, target8
ires4 = SHMEM_INT4_FINC(target4, pe)
ires8 = SHMEM_INT8_FINC(target8, pe)
```
DESCRIPTION

API description

These functions perform a fetch-and-increment operation. The *target* on PE *pe* is increased by one and the function returns the previous contents of *target* as an atomic operation.

Return Values

The contents that had been at the *target* address on the remote PE prior to the increment. The data type of the return value is the same as the *target*.

Notes

None.

EXAMPLES

The following *shmem*_*finc* example is for *C/C++* programs:

```
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19 (mg - 11) ();<br>
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             #include <stdio.h>
             #include <shmem.h>
             int dst;
             int main(void)
             {
                 int me;
                 int old;
                 start_pes(0);
                me = my_pe();old = -1;dst = 22;shmem_barrier_all();
                 if (me == 0)
                    old = shmem_int_finc(\&dst, 1);
                 shmem_barrier_all();
                 printf("%d: old = %d, dst = %d\n", me, old, dst);
                 return 0;
             }
      8.4.5 SHMEM_INC
      Performs an atomic fetch-and-increment operation on a remote data object.
      SYNOPSIS
             C/C++:
             void shmem_int_inc(int *target, int pe);
             void shmem_long_inc(long *target, int pe);
             void shmem_longlong_inc(long long *target, int pe);
             FORTRAN:
             INTEGER pe
             INTEGER(KIND=4) target4
             INTEGER(KIND=8) target8
             CALL SHMEM_INT4_INC(target4, pe)
             CALL SHMEM_INT8_INC(target8, pe)
      DESCRIPTION
             Arguments
                    IN target The remotely accessible integer data object to be updated on the remote
                                                        PE. The type of target should match that implied in the SYNOPSIS
                                                        section.
                    IN pe An integer that indicates the PE number on which target is to be up-
                                                        dated. If you are using Fortran, it must be a default integer value.
             API description
                   These functions perform an atomic increment operation on the target data object on PE.
             Return Values
                   None.
             Notes
                   The term remotely accessible is defined in the Introduction.
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```
EXAMPLES

```
The following shmem_int_inc example is for C/C++ programs:
```

```
#include <stdio.h>
#include <shmem.h>
int dst;
int main(void)
{
  int me;
  start_pes(0);
  me = my_pe();dst = 74;shmem_barrier_all();
  if (me == 0)
      shmem_int_inc(&dst, 1);
   shmem_barrier_all();
  printf("%d: dst = %d\n", me, dst);
   return 0;
}
```
8.4.6 SHMEM**_**FADD

Performs an atomic fetch-and-add operation on a remote data object.

SYNOPSIS

```
C/C++
```

```
int shmem_int_fadd(int *target, int value, int pe);
long shmem_long_fadd(long *target, long value, int pe);
long long shmem_longlong_fadd(long long *target, long long value, int pe);
```
FORTRAN:

```
INTEGER pe
INTEGER(KIND=4) SHMEM_INT4_FADD, ires, target, value
ires = SHMEM_INT4_FADD(target, value, pe)
INTEGER(KIND=8) SHMEM_INT8_FADD, ires, target, value
ires = SHMEM_INT8_FADD(target, value, pe)
```
DESCRIPTION

API description

*shmem*_*fadd* functions perform an atomic fetch-and-add operation. An atomic fetch-and-add operation fetches the old *target* and adds *value* to *target* without the possibility of another atomic operation on the *target* between the time of the fetch and the update. These routines add *value* to *target* on *pe* and return the previous contents of *target* as an atomic operation.

Collective operations are defined as communication or synchronization operations on a group of PEs called *Active set*. The collective operations require all PEs in the *Active set* to simultaneously call the operation. A PE that is not part of the *Active set* calling the collective operations results in an undefined behavior. All collective operations have an *Active set* as an input parameter except *SHMEM*_*BARRIER*_*ALL*. The *SHMEM*_*BARRIER*_*ALL* is called by all PEs of the OpenSHMEM program. 25 26 27 28 29

The *Active set* is defined by the arguments *PE*_*start*, *logPE*_*stride*, and *PE*_*size*. *PE*_*start* is the starting PE number, a log (base 2) of *logPE*_*stride* is the stride between PEs, and *PE*_*size* is the number of PEs participating in the *Active set*. All PEs participating in the collective operations provide the same values for these arguments.

Another argument important to collective operations is *pSync*, which is a symmetric work array. All PEs participating in a collective must pass the same pSync array. On completion of a collective call, the *pSync* is restored to its original contents. The reuse of *pSync* array is allowed for a PE, if all previous collective operations using the *pSync* array is completed by all participating PEs. One can use a synchronization collective operation such as *SHMEM*_*BARRIER* to ensure completion of previous collective operations. The two cases below show the reuse of *pSync* array:

- The *shmem barrier* function allows the same *pSync* array to be used on consecutive calls as long as the active PE set does not change.
- If the same collective function is called multiple times with the same *Active set*, the calls may alternate between two *pSync* arrays. The OpenSHMEM functions guarantee that a first call is completely finished by all PEs by the time processing of a third call begins on any PE.

All collective operations defined in the specification are blocking. The collective operations return on completion. The collective operations defined in the OpenSHMEM specification are:

- *SHMEM*_*BROADCAST* 45
- *SHMEM*_*BARRIER* 47

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*SHMEM*_*BARRIER*_*ALL* 48

*SHMEM*_*COLLECT*

Reduction Operations

8.5.1 SHMEM**_**BARRIER**_**ALL

Registers the arrival of a PE at a barrier and suspends PE execution until all other PEs arrive at the barrier and all local and remote memory updates are completed.

SYNOPSIS

C/C_{++} :

void shmem_barrier_all(**void**); FORTRAN:

CALL SHMEM_BARRIER_ALL

DESCRIPTION

Arguments None.

API description

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 RTRAN:

a sinese particular strength (including mainter all function registers the arrival of a PE at a barrier. Barriers are a fast mecha The *shmem*_*barrier*_*all* function registers the arrival of a PE at a barrier. Barriers are a fast mechanism for synchronizing all PEs at once. This routine causes a PE to suspend execution until all PEs have called *shmem*_*barrier*_*all*. This function must be used with PEs started by *start*_*pes*.

Prior to synchronizing with other PEs, *shmem*_*barrier*_*all* ensures completion of all previously issued memory stores and remote memory updates issued via OpenSHMEM AMOs and RMA routine calls such as *shmem*_*int*_*add* and *shmem*_*put32*.

Return Values

None.

Notes

None.

EXAMPLES

The following *shmem*_*barrier*_*all* example is for *C/C++* programs:

```
#include <stdio.h>
#include <shmem.h>
int x=1010;
```
int main(**void**) {

```
int me, npes;
start_pes(0);
me = my_pe();npes = \text{num}pes();/*put to next PE in a circular fashion*/
shmem_int_p(&x, 4, me+1%npes);
/*synchronize all PEs*/
shmem_barrier_all();
```

```
printf("%d: x = %d\n\pi, me, x);
return 0;
```

```
}
```
8.5.2 SHMEM**_**BARRIER

Performs all operations described in the *shmem*_*barrier*_*all* interface but with respect to a subset of PEs defined by the *Active set*.

SYNOPSIS

C/C++:

void shmem_barrier(**int** PE_start, **int** logPE_stride, **int** PE_size, **long** *pSync);

PE_size

INTEGER pSync(SHMEM_BARRIER_SYNC_SIZE)

CALL SHMEM_BARRIER(PE_start, logPE_stride, PE_size, pSync)

DESCRIPTION

API description

*shmem*_*barrier* is a collective synchronization routine over an *Active set*. Control returns from *shmem*_*barrier* after all PEs in the *Active set* (specified by *PE*_*start*, *logPE*_*stride*, and *PE*_*size*) have called *shmem*_*barrier*. As with all OpenSHMEM collective routines, each of these routines assumes that only PEs in the *Active set* call the routine. If a PE not in the *Active set* calls a OpenSHMEM collective routine, undefined behavior results.

The values of arguments *PE*_*start*, *logPE*_*stride*, and *PE*_*size* must be equal on all PEs in the *Active set*. The same work array must be passed in *pSync* to all PEs in the *Active set*.

*shmem*_*barrier* ensures that all previously issued stores and remote memory updates, including AMOs and RMA operations, done by any of the PEs in the *Active set* are complete before returning.

The same *pSync* array may be reused on consecutive calls to *shmem*_*barrier* if the same active PE set is used.

Return Values

None.

Notes

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#include <stdio.h>

EXAMPLES

{

```
The following barrier example is for C/C++ programs:
```

```
#include <shmem.h>
long pSync[_SHMEM_BARRIER_SYNC_SIZE];
int x = 10101;
```

```
int main(void)
```

```
int me, npes;
```

```
for (int i = 0; i < _SHMEM_BARRIER_SYNC_SIZE; i += 1){
  pSync[i] = _SHMEM_SYNC_VALUE;
}
```

```
start_pes(0);
me = my_pe();
npes = _{num\_pes} ();
```

```
if(me % 2 == 0){
   x = 1000 + me;/*put to next even PE in a circular fashion*/
   shmem_int_p(&x, 4, me+2%npes);
   /*synchronize all even pes*/
   shmem_barrier(0, 1, (npes/2 + npes%2), pSync);
}
printf("%d: x = %d\n\infty", me, x);
return 0;
```

```
}
```
8.5.3 SHMEM**_**BROADCAST

Broadcasts a block of data from one PE to one or more *target* PEs.

SYNOPSIS

```
C/C++:
```

```
void shmem_broadcast32(void *target, const void *source, size_t nlong, int PE_root, int
   PE_start, int logPE_stride, int PE_size, long *pSync);
void shmem_broadcast64(void *target, const void *source, size_t nlong, int PE_root, int
```
PE_start, **int** logPE_stride, **int** PE_size, **long** *pSync);

FORTRAN:

```
Let \lim_{x\to 0} \frac{1}{x^{3(3m+1)}} = \frac{1}{x^{3(3m+1)}}<br>
\frac{1}{x^{3(INTEGER nlong, PE_root, PE_start, logPE_stride, PE_size
INTEGER pSync(SHMEM_BCAST_SYNC_SIZE)
CALL SHMEM_BROADCAST4(target, source, nlong, PE_root, PE_start, logPE_stride, PE_size, pSync)
CALL SHMEM_BROADCAST8(target, source, nlong, PE_root, PE_start, logPE_stride, PE_size, pSync)
CALL SHMEM_BROADCAST32(target, source, nlong, PE_root, PE_start, logPE_stride, PE_size,pSync)
CALL SHMEM_BROADCAST64(target, source, nlong, PE_root, PE_start, logPE_stride, PE_size,pSync)
```
DESCRIPTION

PEs participates in processing of a OpenSHMEM broadcast routine. Be careful to avoid these situations:

If the *pSync* array is initialized at run time, some type of synchronization is needed to ensure that all PEs in the working set have initialized *pSync* before any of them enter a OpenSHMEM routine called with the *pSync* synchronization array. A *pSync* array may be reused on a subsequent OpenSHMEM broadcast routine only if none of the PEs in the *Active set* are still processing a prior OpenSHMEM broadcast routine call that used the same *pSync* array. In general, this can be ensured only by doing some type of synchronization. However, in the special case of OpenSHMEM routines being called with the same *Active set*, you can allocate two *pSync* arrays and alternate between them on successive calls.

EXAMPLES

In the following examples, the call to *shmem*_*broadcast64* copies *source* on PE 4 to *target* on PEs 5, 6, and 7. *C/C++* example:

```
for (i=0; i < \leq SHMEM_BCAST_SYNC_SIZE; i++) {
        pSync[i] = _SHMEM_SYNC_VALUE;
}
shmem_barrier_all(); /* Wait for all PEs to initialize pSync
shmem_broadcast64(target, source, nlong, 0, 4, 0, 4, pSync);
```
Fortran example:

```
INTEGER PSYNC(SHMEM_BCAST_SYNC_SIZE)
INTEGER TARGET, SOURCE, NLONG, PE_ROOT, PE_START,
   & LOGPE_STRIDE, PE_SIZE, PSYNC
COMMON /COM/ TARGET, SOURCE
```
DATA PSYNC /SHMEM_BCAST_SYNC_SIZE*SHMEM_SYNC_VALUE/

CALL SHMEM_BROADCAST64(TARGET, SOURCE, NLONG, 0, 4, 0, 4, PSYNC)

8.5.4 SHMEM**_**COLLECT, SHMEM**_**FCOLLECT

Concatenates blocks of data from multiple PEs to an array in every PE.

SYNOPSIS

```
4+ example:<br>
(1-0) 2 < SOMERLE-CREW_CREW_VALUE}<br>
methodical of \sim Responses and the properties of the state of the s
C/C++void shmem_collect32(void *target, const void *source, size_t nelems, int PE_start, int
     logPE_stride, int PE_size, long *pSync);
void shmem_collect64(void *target, const void *source, size_t nelems, int PE_start, int
    logPE_stride, int PE_size, long *pSync);
void shmem_fcollect32(void *target, const void *source, size_t nelems, int PE_start, int
     logPE_stride, int PE_size, long *pSync);
void shmem_fcollect64(void *target, const void *source, size_t nelems, int PE_start, int
     logPE_stride, int PE_size, long *pSync);
```
FORTRAN:

```
INTEGER nelems
INTEGER PE_start, logPE_stride, PE_size
INTEGER pSync(SHMEM_COLLECT_SYNC_SIZE)
CALL SHMEM_COLLECT4(target, source, nelems, PE_start, logPE_stride, PE_size, pSync)
CALL SHMEM_COLLECT8(target, source, nelems, PE_start, logPE_stride, PE_size, pSync)
CALL SHMEM_COLLECT32(target, source, nelems, PE_start, logPE_stride, PE_size, pSync)
CALL SHMEM_COLLECT64(target, source, nelems, PE_start, logPE_stride, PE_size, pSync)
CALL SHMEM_FCOLLECT4(target, source, nelems, PE_start, logPE_stride, PE_size, pSync)
CALL SHMEM_FCOLLECT8(target, source, nelems, PE_start, logPE_stride, PE_size, pSync)
CALL SHMEM_FCOLLECT32(target, source, nelems, PE_start, logPE_stride, PE_size, pSync)
CALL SHMEM_FCOLLECT64(target, source, nelems, PE_start, logPE_stride, PE_size, pSync)
```
DESCRIPTION

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You must ensure that the *pSync* array is not being updated on any PE in the *Active set* while any of the PEs participate in processing of a OpenSHMEM collective routine. Be careful to avoid these situations: If the *pSync* array is initialized at run time, some type of synchronization is needed to ensure that all PEs in the working set have initialized *pSync* before any of them enter a OpenSHMEM routine called with the *pSync* synchronization array. A *pSync* array can be reused on a subsequent OpenSHMEM collective routine only if none of the PEs in the *Active set* are still processing a prior OpenSHMEM collective routine call that used the same *pSync* array. In general, this may be ensured only by doing some type of synchronization. However, in the special case of OpenSHMEM routines being called with the same *Active set*, you can allocate two *pSync* arrays and alternate between them on successive calls.

The collective routines operate on active PE sets that have a non-power-of-two *PE*_*size* with some performance degradation. They operate with no performance degradation when *nelems* is a non-power-of-two value.

EXAMPLES

The following *shmem*_*collec*t example is for *C/C++* programs:

```
for (i=0; i < _SHMEM_COLLECT_SYNC_SIZE; i++) {
         pSync[i] = _SHMEM_SYNC_VALUE;
}
shmem_barrier_all(); /* Wait for all PEs to initialize pSync
shmem_collect32(target, source, 64, pe_start, logPE_stride,
  pe_size, pSync);
```
The following *SHMEM*_*COLLECT* example is for *Fortran* programs:

INTEGER PSYNC(SHMEM_COLLECT_SYNC_SIZE) **DATA** PSYNC /SHMEM_COLLECT_SYNC_SIZE*SHMEM_SYNC_VALUE/

CALL SHMEM_COLLECT4(TARGET, SOURCE, 64, PE_START, LOGPE_STRIDE, & PE_SIZE, PSYNC)

8.5.5 SHMEM**_**REDUCTIONS

Performs a logical operations across a set of PEs.

SYNOPSIS


```
FELIENT, int logFE_skiele, int FE_skie, long dochle -yKK, long (pSyne);<br>
FELIENT, interpretise and the particular interpretise and the syne of the reduces of the reduces of the syne of the content, interpretise and the re
           void shmem_short_max_to_all(short *target, short *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, short *pWrk, long *pSync);
           void shmem_double_min_to_all(double *target, double *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, double *pWrk, long *pSync);
           void shmem_float_min_to_all(float *target, float *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, float *pWrk, long *pSync);
           void shmem_int_min_to_all(int *target, int *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, int *pWrk, long *pSync);
           void shmem_long_min_to_all(long *target, long *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, long *pWrk, long *pSync);
           void shmem_longdouble_min_to_all(long double *target, long double *source, int nreduce, int
               PE_start, int logPE_stride, int PE_size, long double *pWrk, long *pSync);
           void shmem_longlong_min_to_all(long long *target, long long *source, int nreduce, int
               PE_start, int logPE_stride, int PE_size, long long *pWrk, long *pSync);
           void shmem_short_min_to_all(short *target, short *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, short *pWrk, long *pSync);
           void shmem_complexd_sum_to_all(double complex *target, double complex *source, int nreduce,
               int PE_start, int logPE_stride, int PE_size, double complex *pWrk, long *pSync);
           void shmem_complexf_sum_to_all(float complex *target, float complex *source, int nreduce, int
                PE_start, int logPE_stride, int PE_size, float complex *pWrk, long *pSync);
           void shmem_double_sum_to_all(double *target, double *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, double *pWrk, long *pSync);
           void shmem_float_sum_to_all(float *target, float *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, float *pWrk, long *pSync);
           void shmem_int_sum_to_all(int *target, int *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, int *pWrk, long *pSync);
           void shmem_long_sum_to_all(long *target, long *source, int nreduce, int PE_start, int
               logPE_stride,int PE_size, long *pWrk, long *pSync);
           void shmem_longdouble_sum_to_all(long double *target, long double *source, int nreduce, int
               PE_start, int logPE_stride, int PE_size, long double *pWrk, long *pSync);
           void shmem_longlong_sum_to_all(long long *target, long long *source, int nreduce, int
               PE_start, int logPE_stride, int PE_size, long long *pWrk, long *pSync);
           void shmem_short_sum_to_all(short *target, short *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, short *pWrk, long *pSync);
           void shmem_complexd_prod_to_all(double complex *target, double complex *source, int nreduce,
               int PE_start, int logPE_stride, int PE_size, double complex *pWrk, long *pSync);
           void shmem_complexf_prod_to_all(float complex *target, float complex *source, int nreduce,
               int PE_start, int logPE_stride, int PE_size, float complex *pWrk, long *pSync);
           void shmem_double_prod_to_all(double *target, double *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, double *pWrk, long *pSync);
           void shmem_float_prod_to_all(float *target, float *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, float *pWrk, long *pSync);
           void shmem_int_prod_to_all(int *target, int *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, int *pWrk, long *pSync);
           void shmem_long_prod_to_all(long *target, long *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, long *pWrk, long *pSync);
           void shmem_longdouble_prod_to_all(long double *target, long double *source, int nreduce, int
               PE_start, int logPE_stride, int PE_size, long double *pWrk, long *pSync);
           void shmem_longlong_prod_to_all(long long *target, long long *source, int nreduce, int
               PE_start, int logPE_stride, int PE_size, long long *pWrk, long *pSync);
           void shmem_short_prod_to_all(short *target, short *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, short *pWrk, long *pSync);
           void shmem_int_or_to_all(int *target, int *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, int *pWrk, long *pSync);
           void shmem_long_or_to_all(long *target, long *source, int nreduce, int PE_start, int
               logPE_stride, int PE_size, long *pWrk, long *pSync);
           void shmem_longlong_or_to_all(long long *target, long long *source, int nreduce, int PE_start
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```
pSync)

RTRAN:

2010-11-2000_TO_ALLitarget, source, produce, PE_start, logPE_acride, PE_stare, pRt.

2020-11

2 **void** shmem_short_or_to_all(**short** *target, **short** *source, **int** nreduce, **int** PE_start, **int** logPE_stride, **int** PE_size, **short** *pWrk, **long** *pSync); **void** shmem_int_xor_to_all(**int** *target, **int** *source, **int** nreduce, **int** PE_start, **int** logPE_stride, **int** PE_size, **int** *pWrk, **long** *pSync); **void** shmem_long_xor_to_all(**long** *target, **long** *source, **int** nreduce, **int** PE_start, **int** logPE_stride, **int** PE_size, **long** *pWrk, **long** *pSync); **void** shmem_longlong_xor_to_all(**long long** *target, **long long** *source, **int** nreduce, **int** PE_start, **int** logPE_stride, **int** PE_size, **long long** *pWrk, **long** *pSync); **void** shmem_short_xor_to_all(**short** *target, **short** *source, **int** nreduce, **int** PE_start, **int** logPE_stride, **int** PE_size, **short** *pWrk, **long** *pSync); FORTRAN: **CALL** SHMEM_INT4_AND_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_INT8_AND_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_INT4_MAX_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_INT8_MAX_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_REAL4_MAX_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_REAL8_MAX_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_REAL16_MAX_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_INT4_MIN_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_INT8_MIN_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_REAL4_MIN_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_REAL8_MIN_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) CALL SHMEM_REAL16_MIN_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_COMP4_SUM_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_COMP8_SUM_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_INT4_SUM_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_INT8_SUM_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_REAL4_SUM_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_REAL8_SUM_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_REAL16_SUM_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, pSync) **CALL** SHMEM_COMP4_PROD_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk, 17 18 19 20 21 22 23 24 25 26 27 28 29 31 32 33 34 35 36 37 39 41 42 43 47

, **int** logPE_stride, **int** PE_size, **long long** *pWrk, **long** *pSync);

CALL SHMEM_COMP8_PROD_TO_ALL(target, source, nreduce, PE_start, logPE_stride, PE_size, pWrk,

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DESCRIPTION

API description

OpenSHMEM reduction routines compute one or more reductions across symmetric arrays on multiple virtual PEs. A reduction performs an associative binary operation across a set of values.

The *nreduce* argument determines the number of separate reductions to perform. The *source* array on all PEs in the *Active set* provides one element for each reduction. The results of the reductions are placed in the *target* array on all PEs in the *Active set*. The *Active set* is defined by the *PE*_*start*, *logPE*_*stride*, *PE*_*size* triplet.

The *source* and *target* arrays may be the same array, but they may not be overlapping arrays.

As with all OpenSHMEM collective routines, each of these routines assumes that only PEs in the *Active set* call the routine. If a PE not in the *Active set* calls a OpenSHMEM collective routine, undefined behavior results.

The values of arguments *nreduce*, *PE*_*start*, *logPE*_*stride*, and *PE*_*size* must be equal on all PEs in the *Active set*. The same *target* and *source* arrays, and the same *pWrk* and *pSync* work arrays, must be passed to all PEs in the *Active set*.

Before any PE calls a reduction routine, you must ensure that the following conditions exist (synchronization via a *barrier* or some other method is often needed to ensure this): The *pWrk* and *pSync* arrays on all PEs in the *Active set* are not still in use from a prior call to a collective OpenSHMEM routine. The *target* array on all PEs in the *Active set* is ready to accept the results of the *reduction*.

Upon return from a reduction routine, the following are true for the local PE: The *target* array is updated. The values in the *pSync* array are restored to the original values.

When calling from *Fortran*, the *target* date types are as follows:

Return Values

None.

Notes

All OpenSHMEM reduction routines reset the values in *pSync* before they return, so a particular *pSync* buffer need only be initialized the first time it is used.

shmen_int8_xor_to_all

shmen_int8_xor_to_all

shmen_int8_xor_to_all

shmen_int8_xor_to_all

shmen_intal Lateger, with an element size of 8 bytes.

shmen_reall4_xor_to_all

Real, with an element size of 8 bytes.

shmen_rea You must ensure that the *pSync* array is not being updated on any PE in the *Active set* while any of the PEs participate in processing of a OpenSHMEM reduction routine. Be careful to avoid the following situations: If the *pSync* array is initialized at run time, some type of synchronization is needed to ensure that all PEs in the working set have initialized *pSync* before any of them enter an OpenSHMEM routine called with the *pSync* synchronization array. A *pSync* or *pWrk* array can be reused in a subsequent reduction routine call only if none of the PEs in the *Active set* are still processing a prior reduction routine call that used the same *pSync* or *pWrk* arrays. In general, this can be assured only by doing some type of synchronization. However, in the special case of reduction routines being called with the same *Active set*, you can allocate two *pSync* and *pWrk* arrays and alternate between them on successive calls.

EXAMPLES

This *Fortran* example statically initializes the *pSync* array and finds the logical *AND* of the integer variable *FOO* across all even PEs.

```
INCLUDE "shmem.fh"
INTEGER PSYNC(SHMEM_REDUCE_SYNC_SIZE)
DATA PSYNC /SHMEM_REDUCE_SYNC_SIZE*SHMEM_SYNC_VALUE/
PARAMETER (NR=1)
REAL PWRK(MAX(NR/2+1,SHMEM_REDUCE_MIN_WRKDATA_SIZE))
INTEGER FOO, FOOAND
COMMON /COM/ FOO, FOOAND, PWRK
INTRINSIC MY_PE
IF ( MOD(MY_PE(),2) .EQ. 0) THEN
       CALL SHMEM_INT8_AND_TO_ALL(FOOAND, FOO, NR, 0, 1, N$PES/2,
& PWRK, PSYNC)
       PRINT*,'Result on PE ',MY_PE(),' is ',FOOAND
ENDIF
```
This *Fortran* example statically initializes the *pSync* array and finds the *maximum* value of real variable *FOO* across all even PEs. 46

```
INCLUDE "shmem.fh"
           INTEGER PSYNC(SHMEM_REDUCE_SYNC_SIZE)
           DATA PSYNC /SHMEM_REDUCE_SYNC_SIZE*SHMEM_SYNC_VALUE/
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```
8. OPENSHMEM LIBRARY API 45

```
PARAMETER (NR=1)
REAL FOO, FOOMAX, PWRK(MAX(NR/2+1,SHMEM_REDUCE_MIN_WRKDATA_SIZE))
COMMON /COM/ FOO, FOOMAX, PWRK
INTRINSIC MY_PE
IF ( MOD(MY_PE(),2) .EQ. 0) THEN
      CALL SHMEM_REAL8_MAX_TO_ALL(FOOMAX, FOO, NR, 0, 1, N$PES/2,
& PWRK, PSYNC)
      PRINT*,'Result on PE ',MY_PE(),' is ',FOOMAX
ENDIF
```
This *Fortran* example statically initializes the *pSync* array and finds the *minimum* value of real variable *FOO* across all the even PEs.

```
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A PEOPE / SINGLE RECOCT DEVICATION SYNC, VALUE /<br>
EVALUATION CONTEST<br>
EVALUATIONS:
INCLUDE "shmem.fh"
INTEGER PSYNC(SHMEM_REDUCE_SYNC_SIZE)
DATA PSYNC /SHMEM_REDUCE_SYNC_SIZE*SHMEM_SYNC_VALUE/
PARAMETER (NR=1)
REAL FOO, FOOMIN, PWRK(MAX(NR/2+1,SHMEM_REDUCE_MIN_WRKDATA_SIZE))
COMMON /COM/ FOO, FOOMIN, PWRK
INTRINSIC MY_PE
IF ( MOD(MY_PE(),2) .EQ. 0) THEN
       CALL SHMEM_REAL8_MIN_TO_ALL(FOOMIN, FOO, NR, 0, 1, N$PES/2,
& PWRK, PSYNC)
        PRINT*,'Result on PE ',MY_PE(),' is ',FOOMIN
ENDIF
This Fortran example statically initializes the pSync array and finds the sum of the real variable FOO across all
even PEs.
INCLUDE "shmem.fh"
INTEGER PSYNC(SHMEM_REDUCE_SYNC_SIZE)
DATA PSYNC /SHMEM_REDUCE_SYNC_SIZE*SHMEM_SYNC_VALUE/
PARAMETER (NR=1)
REAL FOO, FOOSUM, PWRK(MAX(NR/2+1,SHMEM_REDUCE_MIN_WRKDATA_SIZE))
COMMON /COM/ FOO, FOOSUM, PWRK
INTRINSIC MY_PE
IF ( MOD(MYPE(), 2) .EQ. 0) THEN
        CALL SHMEM_INT4_SUM_TO_ALL(FOOSUM, FOO, NR, 0, 1, N$PES/2,
& PWRK, PSYNC)
        PRINT*,'Result on PE ',MY_PE(),' is ',FOOSUM
ENDIF
This Fortran example statically initializes the pSync array and finds the product of the real variable FOO across
all the even PEs.
INCLUDE "shmem.fh"
INTEGER PSYNC(SHMEM_REDUCE_SYNC_SIZE)
DATA PSYNC /SHMEM_REDUCE_SYNC_SIZE*SHMEM_SYNC_VALUE/
PARAMETER (NR=1)
REAL FOO, FOOPROD, PWRK(MAX(NR/2+1,SHMEM_REDUCE_MIN_WRKDATA_SIZE))
COMMON /COM/ FOO, FOOPROD, PWRK
INTRINSIC MY_PE
IF ( MOD(MY_PEC(), 2) .EQ. 0) THEN
        CALL SHMEM_COMP8_PROD_TO_ALL(FOOPROD, FOO, NR, 0, 1, N$PES/2,
   PWRK, PSYNC)
         PRINT*,'Result on PE ',MY_PE(),' is ',FOOPROD
ENDIF
```
This *Fortran* example statically initializes the *pSync* array and finds the logical *OR* of the integer variable *FOO* across all even PEs.

```
INCLUDE "mpp/shmem.fh"
           INTEGER PSYNC(SHMEM_REDUCE_SYNC_SIZE)
           DATA PSYNC /SHMEM_REDUCE_SYNC_SIZE*SHMEM_SYNC_VALUE/
           PARAMETER (NR=1)
           REAL PWRK(MAX(NR/2+1,SHMEM_REDUCE_MIN_WRKDATA_SIZE))
           INTEGER FOO, FOOOR
           COMMON /COM/ FOO, FOOOR, PWRK
           INTRINSIC MY_PE
           IF ( MOD(MY_PE(),2) .EQ. 0) THEN
                    CALL SHMEM_INT8_OR_TO_ALL(FOOOR, FOO, NR, 0, 1, N$PES/2,
           & PWRK, PSYNC)
                    PRINT*,'Result on PE ',MY_PE(),' is ',FOOOR
           ENDIF
           This Fortran example statically initializes the pSync array and computes the exclusive XOR of variable FOO
           across all even PEs.
           INCLUDE "mpp/shmem.fh"
           INTEGER PSYNC(SHMEM_REDUCE_SYNC_SIZE)
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```
PRINT, Result on PE<sup>2</sup>, NY_PE(),<sup>2</sup> is <sup>2</sup>, NY_OOR<br>
K Fortran example statically initializes the pSync array and computes the exclusive XOR of variable F(<br>
K Fortran example statically initializes the pSync 
DATA PSYNC /SHMEM_REDUCE_SYNC_SIZE*SHMEM_SYNC_VALUE/
PARAMETER (NR=1)
REAL FOO, FOOXOR, PWRK(MAX(NR/2+1,SHMEM_REDUCE_MIN_WRKDATA_SIZE))
COMMON /COM/ FOO, FOOXOR, PWRK
INTRINSIC MY_PE
IF ( MOD(MY_PE(),2) .EQ. 0) THEN
         CALL SHMEM_REAL8_XOR_TO_ALL(FOOXOR, FOO, NR, 0, 1, N$PES/2,
& PWRK, PSYNC)
         PRINT*,'Result on PE ',MY_PE(),' is ',FOOXOR
```
8.6 Point-to-point synchronization functions

The following section discusses OpenSHMEM API that provides a mechanism for synchronization between two PEs based on the value of a symmetric data object.

8.6.1 SHMEM**_**WAIT

ENDIF

Wait for a variable on the local PE to change.

SYNOPSIS

```
C/C++:
           void shmem_int_wait(int *var, int value);
           void shmem_int_wait_until(int *var, int cond, int value);
           void shmem_long_wait(long *var, long value);
           void shmem_long_wait_until(long *var, int cond, long value);
           void shmem_longlong_wait(long long *var, long long value);
           void shmem_longlong_wait_until(long long *var, int cond, long long value);
           void shmem_short_wait(short *var, short value);
           void shmem_short_wait_until(short *var, int cond, short value);
           void shmem_wait(long *ivar, long cmp_value);
           void shmem_wait_until(long *ivar, int cmp, long value);
           FORTRAN:
           CALL SHMEM_INT4_WAIT(ivar, cmp_value)
           CALL SHMEM_INT4_WAIT_UNTIL(ivar, cmp, cmp_value)
           CALL SHMEM_INT8_WAIT(ivar, cmp_value)
           CALL SHMEM_INT8_WAIT_UNTIL(ivar, cmp, cmp_value)
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```
8. OPENSHMEM LIBRARY API 47

```
CALL SHMEM_WAIT(ivar, cmp_value)
CALL SHMEM_WAIT_UNTIL(ivar, cmp, cmp_value)
```
DESCRIPTION

API description

shmem_wait and shmem_wait_until wait for ivar to be changed by a remote write or atomic swap issued by a different processor. These routines can be used for point-to-point directed synchronization. A call to shmem_wait does not return until some other processor writes a value, not equal to cmp_value, into ivar on the waiting processor. A call to shmem_wait_until does not return until some other processor changes ivar to satisfy the condition implied by cmp and cmp_value. This mechanism is useful when a processor needs to tell another processor that it has completed some action. The shmem_wait routines return when ivar is no longer equal to cmp_value. The shmem_wait_until routines return when the compare condition is true. The compare condition is defined by the ivar argument compared with the cmp_value using the comparison operator, cmp.

If you are using *Fortran*, ivar must be a specific sized integer type according to the function being called, as follows:

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8.7.1 SHMEM**_**FENCE

Assures ordering of delivery of *Put*, AMOs, and store operations to symmetric data objects. 41

```
SYNOPSIS
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```


DESCRIPTION

Arguments

None.

API description

This function assures ordering of delivery of *Put*, AMOs, and store operations to symmetric data objects. All *Put*, AMOs, and store operations to symmetric data objects issued to a particular remote PE prior to the call to *shmem*_*fence* are guaranteed to be ordered to be delivered before any subsequent *Put*, AMOs, and store operations to symmetric data objects to the same PE.

Return Values

None.

Notes

*shmem*_*fence* only provides per-PE ordering guarantees and does not guarantee completion of delivery. There is a subtle difference between *shmem*_*fence* and *shmem*_*quiet*, in that, that *shmem*_*quiet* guarantees completion of *Put*, AMOs, and store operations to symmetric data objects which makes the updates visible to all other PEs.

The *shmem*_*quiet* function should be called if completion of PUT, AMOs, and store operations to symmetric data objects is desired when multiple remote PEs are involved.

EXAMPLES

The following *shmem*_*fence* example is for *C/C++* programs:

```
#include <stdio.h>
#include <shmem.h>
```

```
long target [10] = \{0\};
int targ = 0;
```

```
int main(void)
```

```
urn Values<br>
Some.<br>
Somewhere figure only provides per-PE ordering guarantees and does not guarantee completion of deliver<br>
There is a subtle difference between shmem_ferce and shmem<sup>f</sup>aguick, in that, that shmem_qu
{
  long source[10] = { 1, 2, 3, 4, 5, 6, 7, 8, 9, 10 };
 int src = 99;
 start_pes(0);
  if (_my_pe() == 0) {
    shmem_long_put(target, source, 10, 1); /*put1*/
    shmem_long_put(target, source, 10, 2); /*put2*/
    shmem_fence();
    shmem_int_put(&targ, &src, 1, 1); /*put3*/
    shmem_int_put(&targ, &src, 1, 2); /*put4*/
  }
  shmem_barrier_all(); /* sync sender and receiver */
  printf("target[0] on PE %d is %d\n", _my_pe(), target[0]);
  return 1;
```

```
}
```
Put1 will be ordered to be delivered before *put3* and *put2* will be ordered to be delivered before *put4*.

8.7.2 SHMEM**_**QUIET

Waits for completion of all outstanding *Put*, AMOs and store operations to symmetric data objects issued by a PE.

SYNOPSIS

 $C/C++$:

void shmem_quiet(**void**);

8.7.3 Synchronization and Communication Ordering in OpenSHMEM

When using the OpenSHMEM API, synchronization, ordering, and completion of communication become critical. The updates via *Put* operations, AMOs and store operations on symmetric data cannot be guaranteed until some form of synchronization or ordering is introduced by the application programmer. The table below gives the different synchronization and ordering choices, and the situations where they may be useful.

8.8 Distributed Locking Operations

The following section discusses OpenSHMEM locks as a mechanism to provide mutual exclusion. Three operations are available for distributed locking, *set, test* and *clear*.

8.8.1 SHMEM**_**LOCK

Releases, locks, and tests a mutual exclusion memory lock.

```
A sharen_clear_lock (long -lock);<br>
sheen_clear_lock (long -lock);<br>
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       SYNOPSIS
             C/C++:
             void shmem_clear_lock(long *lock);
             void shmem_set_lock(long *lock);
             int shmem_test_lock(long *lock);
             FORTRAN:
             INTEGER lock, SHMEM_TEST_LOCK
             CALL SHMEM_CLEAR_LOCK(lock)
             CALL SHMEM_SET_LOCK(lock)
             I = SHMEM_TEST_LOCK(lock)
      DESCRIPTION
             Arguments
                    IN lock A symmetric data object that is a scalar variable or an array of length 1.
                                                         This data object must be set to 0 on all PEs prior to the first use. lock
                                                         must be of type long. If you are using Fortran, it must be of default
                                                         kind.
             API description
                   The shmem_set_lock routine sets a mutual exclusion lock after waiting for the lock to be freed by any other
                   PE currently holding the lock. Waiting PEs are assured of getting the lock in a first-come, first-served man-
                   ner. The shmem_clear_lock routine releases a lock previously set by shmem_set_lock after ensuring that
                   all local and remote stores initiated in the critical region are complete. The shmem_test_lock function sets
                   a mutual exclusion lock only if it is currently cleared. By using this function, a PE can avoid blocking on
                   a set lock. If the lock is currently set, the function returns without waiting. These routines are appropriate
                   for protecting a critical region from simultaneous update by multiple PEs.
             Return Values
                   The shmem_test_lock function returns 0 if the lock was originally cleared and this call was able to set the
                   lock. A value of 1 is returned if the lock had been set and the call returned without waiting to set the lock.
             Notes
                   The term symmetric data object is defined in Introduction. The lock variable should always be initialized
                   to zero and accessed only by the OpenSHMEM locking API. Changing the value of the lock variable by
                   other means without using the OpenSHMEM API, can lead to undefined behavior.
      EXAMPLES
             The following simple example uses shmem_lock in a C program.
             #include <stdio.h>
             #include <shmem.h>
             long L = 0;
             int main(int argc, char **argv)
             {
                 int me, slp;
                 start_pes(0);
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```

```
me = my_pe();slp = 1;shmem_barrier_all();
  if (me == 1)
     sleep (3);
  shmem_set_lock(&L);
  printf("%d: sleeping %d second%s...\n", me, slp, slp == 1 ? "" : "s");
  sleep(slp);
  printf("%d: sleeping...done\n", me);
  shmem_clear_lock(&L);
  shmem_barrier_all();
  return 0;
}
```
8.9 Deprecated API

All of these operations are deprecated and are provided for backwards compatibility. Implementations must include all items in this section and the operations should function properly, while notifying the user about deprecation of the functionality.

8.9.1 SHMEM**_**CACHE

Controls data cache utilities.

SYNOPSIS

DESCRIPTION

Arguments

IN *target* A data object that is local to the PE. *target* can be of any noncharacter type. If you are using *Fortran*, it can be of any kind.

API description

*shmem*_*set*_*cache*_*inv* enables automatic cache coherency mode. *shmem*_*set*_*cache*_*line*_*inv* enables automatic cache coherency mode for the cache line associated with the address of *target* only. *shmem*_*clear*_*cache*_*inv* disables automatic cache coherency mode previously enabled by *shmem*_*set*_*cache*_*inv* 44 or *shmem*_*set*_*cache*_*line*_*inv*. *shmem*_*udcflush* makes the entire user data cache coherent. *shmem*_*udcflush*_*line* makes coherent the cache line that corresponds with the address specified by *target*. 41 42 43 45 46 47

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Annex A

Writing OpenSHMEM Programs

Incorporating OpenSHMEM into Programs

In this section we describe how to write a "Hello World" OpenSHMEM program. To write a "Hello World" OpenSHMEM program we need to

- Add the include file shmem.h (for *C*) or shmem.fh (for *Fortran*).
- Add the initialization call *start_pes*, (line 9) use single integer argument, 0, which is ignored ¹.
- Use OpenSHMEM calls to query the the total number of PEs (line 10) and PE id (line 11).
- There is no explicit finalize call, either a return from main() (line 13) or an explicit exit() acts as an implicit OpenSHMEM finalization.
- In OpenSHMEM the order in which lines appear in the output is not fixed as PEs execute asynchronously in parallel.

```
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DRAFTING OPEN SHAREM into Programs<br>
into we describe how to write a "Hello World" OpenSHMEM program. To write a "Hello World" OpenSH<br>
the include file shmem.h (for C) or shmem.fh (f
 1 #include <stdio.h>
2 #include <shmem.h> /* The shmem header file */
3
4 int
5 main (int argc, char *argv[])
6 {
7 int nprocs, me;
8
9 start_pes (0);
10 nprocs = shmem_n_pes ();
11 me = shmem_my_pe ();
12 printf ("Hello from %d of %d\n", me, nprocs);
13 return 0;
14 }
```
Listing A.1: Expected Output (4 processors)

 Hello from 0 of 4 Hello from 2 of 4 Hello from 3 of 4 Hello from 1 of 4

OpenSHMEM also has a *Fortran* API, so for completeness we will now give the same program written in *Fortran*, in listing [A:](#page-61-0)

¹The unused argument is for compatibility with older SHMEM implementations.

```
Format (The Lie Framel, 13X, 14, 13X, f(x, y, z), 15, 15)<br>
Expansion 184110<br>
Expansion 18412<br>
Expansi
       1 program hello
       2
       3 include 'shmem.fh'
       4 integer :: shmem_my_pe, shmem_n_pes
       5
       6 integer :: npes, me
       7
       8 call start_pes (0)
             9 npes = shmem_n_pes ()
      10 me = shmem_my_pe ()
      11
      12 write (*, 1000) me, npes
      13
      14 1000 format ('Hello from', 1X, I4, 1X, 'of', 1X, I4)
      \frac{15}{16}16 end program hello
                                                 Listing A.2: Expected Output (4 processors)
       1 Hello from 0 of 4
       2 Hello from 2 of 4
       3 Hello from 3 of 4
       4 Hello from 1 of 4
          The following example shows a more complex OpenSHMEM program that illustrates the use of symmetric data
       objects. Note the declaration of the static short target array and its use as the remote destination in OpenSHMEM short
       Put. The use of the static keyword results in the target array being symmetric on PE 0 and PE 1. Each PE is able to
       transfer data to the target array by simply specifying the local address of the symmetric data object which is to receive
       the data. This aids programmability, as the address of the target need not be exchanged with the active side (PE 0) prior
       to the RMA (Remote Memory Access) operation. Conversely, the declaration of the short source array is asymmetric.
       Because the Put handles the references to the source array only on the active (local) side, the asymmetric source object
      is handled correctly.
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```
For (i-1), i \le 37/2; i \ne 1<br>
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\ne 2.02.93444761 = 2002.000 and 2.0 \ne/<br>
\ne 0.04.7544677 working the starget on each PE \ne/<br>
\{x, y, z\} for (i-1), i \le 37/2; i \le 10, \{x, y\} 
1 #include <shmem.h><br>2 #define SIZE 16
    2 #define SIZE 16
3 int
    4 main(int argc, char* argv[])
5 {
6 short source[SIZE];
7 static short target [SIZE];<br>8 int i;
8 int i;<br>9 int num
9 int num_pe = \text{num\_pes}();<br>10 start pes(0):
10 start_pes(0);<br>11 if (_my_pe()
              if (my_pe() == 0) {12 \frac{1}{\sqrt{\frac{1}{\pi}}}\times \frac{1}{\pi} initialize array \frac{x}{\pi} 13
                        for(i = 0; i < SIZE; i++)
14 source[i] = i;15 /* local, not symmetric */<br>16 /* static makes it symmetr.
16 \frac{1}{4} static makes it symmetric \frac{1}{4}<br>17 \frac{1}{4} put "size" words into target
17 /* put "size" words into target on each PE */
                        for(i = 1; i < num_pe; i++)
19 shmem_short_put(target, source, SIZE, i);
\frac{20}{21}21 shmem_barrier_all(); /* sync sender and receiver */<br>22 if (_my_pe() != 0) {
22 if (_my_pe() != 0) {
                        23 printf("target on PE %d is \t", _my_pe());
24 for(i = 0; i < SIZE; i++)
25 printf("\n");<br>26 printf("\n");
                        printf(\sqrt[n]{n});
27 }
28 shmem_barrier_all(); /* sync before exiting */<br>29 return 0:
              29 return 0;
30 }
```
Listing A.3: Expected Output (4 processors)

Annex B

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Compiling and Running Applications

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the OpenSHMEM specification is silent regarding how OpenSHMEM programs are compiled, linked a

to compile and launch applications.

Interval to compile and launch applications.
 As of now the OpenSHMEM specification is silent regarding how OpenSHMEM programs are compiled, linked and run. This section shows some examples of how wrapper programs are utilized in the OpenSHMEM Reference Implementation to compile and launch applications.

1 Compilation

Applications written in *C*

The OpenSHMEM Reference Implementation provides a wrapper program named oshcc, to aid in the compilation of *C* applications, the wrapper could be called as follows:

oshcc <compiler options> -o myprogram myprogram.c

Where the \langle compiler options \rangle are options understood by the underlying *C* compiler. 24 25

Applications written in *C++*

- The OpenSHMEM Reference Implementation provides a wrapper program named oshCC, to aid in the compilation of *C++* applications, the wrapper could be called as follows: 28 29
- oshCC <compiler options> -o myprogram myprogram.cpp 30
- Where the \langle compiler options \rangle are options understood by the underlying *C++* compiler called by **oshCC**. 31
- Applications written in *Fortran* 33
- The OpenSHMEM Reference Implementation provides a wrapper program named oshfort, to aid in the compilation of *Fortran* applications, the wrapper could be called as follows: 34 35
- oshfort <compiler options> -o myprogram myprogram.f 36
- Where the \langle compiler options) are options understood by the underlying *Fortran* compiler called by **oshfort**. 37 38

2 Running Applications

The OpenSHMEM Reference Implementation provides a wrapper program named oshrun, to launch OpenSHMEM applications, the wrapper could be called as follows: 41 42

oshrun <additional options> -np <#> <program> <program arguments> 43 44

Annex C

Undefined Behavior in OpenSHMEM

The specification provides guidelines to the expected behavior of various library routines. In cases where routines are improperly used or the input is not in accordance with the specification, undefined behavior may be observed. Depending on the implementation there are many interpretations of undefined behavior.

Annex D

Interoperability with other Programming Models

1 MPI Interoperability

OpenSHMEM functions can be used in conjunction with MPI functions in the same application. For example, on SGI systems, programs that use both MPI and OpenSHMEM functions call *MPI*_*Init* and *MPI*_*Finalize* but omit the call to the *start*_*pes* function. OpenSHMEM PE numbers are equal to the MPI rank within the *MPI*_*COMM*_*WORLD* environment variable. Note that this precludes use of OpenSHMEM functions between processes in different *MPI*_*COMM*_*WORLD*s. MPI processes started using the *MPI*_*Comm*_*spawn* function, for example, cannot use OpenSHMEM functions to communicate with their parent MPI processes.

HERE ANDERE AND INTERNATION CONTROLLER THE USE CARRED TO THE CONTROLLER THE COMPARTMENT (ESS) THE COMPARTMENT INTERNATION CONTROLLER THE CARRED TO THE CAPS (CONTROLLER TO THE CAPS (CONTROLLER TO THE CAPS (CONTROLLER TO T On SGI systems MPI jobs that use TCP/sockets for inter-host communication, OpenSHMEM functions can be used to communicate with processes running on the same host. The *shmem*_*pe*_*accessible* function can be used to determine if a remote PE is accessible via OpenSHMEM communication from the local PE. When running an MPI application involving multiple executable files, OpenSHMEM functions can be used to communicate with processes running from the same or different executable files, provided that the communication is limited to symmetric data objects. On these systems, static memory, such as a *Fortran* common block or *C* global variable, is symmetric between processes running from the same executable file, but is not symmetric between processes running from different executable files. Data allocated from the symmetric heap (*shmalloc* or *shpalloc*) is symmetric across the same or different executable files. The function *shmem*_*addr*_*accessible* can be used to determine if a local address is accessible via OpenSHMEM communication from a remote PE.

Another important feature of these systems is that the *shmem*_*pe*_*accessible* function returns *TRUE* only if the remote PE is a process running from the same executable file as the local PE, indicating that full OpenSHMEM support (static memory and symmetric heap) is available. When using OpenSHMEM functions within an MPI program, the use of MPI memory placement environment variables is required when using non-default memory placement options.

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Annex E

History of OpenSHMEM

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- A SHMEM Timeline
	- Cray SHMEM
		- * SHMEM first introduced by Cray Research Inc. in 1993 for Cray T3D
		- * Cray is acquired by SGI in 1996
		- * Cray is acquired by Tera in 2000 (MTA)
		- * Platforms: Cray T3D, T3E, C90, J90, SV1, SV2, X1, X2, XE, XMT, XT
	- SGI SHMEM
		- * SGI purchases Cray Research Inc. and SHMEM was integrated into SGI's Message Passing Toolkit (MPT)

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- * SGI currently owns the rights to SHMEM and OpenSHMEM
- * Platforms: Origin, Altix 4700, Altix XE, Altix ICE, Altix UV
- * SGI was purchased by Rackable Systems in 2009
- * SGI and Open Source Software Solutions, Inc. (OSSS) signed a SHMEM trademark licensing agreement, in 2010
- Other Implementations
	- * Quadrics (Vega UK, Ltd.)
	- * Hewlett Packard
	- * GPSHMEM
	- * IBM
	- * QLogic
	- * Mellanox
	- * University of Florida
- OpenSHMEM Implementations
	- SGI OpenSHMEM
	- University of Houston OpenSHMEM Reference Implementation
	- Mellanox ScalableSHMEM
	- Portals-SHMEM
- Implementations that support OpenSHMEM- *Pending verification*
	- IBM OpenSHMEM

Annex F

Changes to this Document

1 Version 1.1

EXECTS TO THIS [D](#page-52-1)OCUMENTE

SiON 1.1

on summarizes the changes from the OpenSHMEM specification Version 1.0 to the Version 7.1. A mathis version is that it provides an accurate description of OpenSHMEM interfaces so that This section summarizes the changes from the OpenSHMEM specification Version 1.0 to the Version 1.1. A major change in this version is that it provides an accurate description of OpenSHMEM interfaces so that they are in agreement with the SGI specification. This version also explains OpenSHMEM 's programming, memory, and execution model. The document was throughly changed to improve the readability of specification and usability of interfaces. The code examples were added to demonstrate the usability of API. Additionally, diagrams were added to help understand the subtle semantic differences of various operations.

The following list describes the specific changes in 1.1:

- Clarifications on the completion semantics of memory synchronization interfaces. See Section 8.7.
- Clarification about completion semantics of memory load and store operations in context of *shmem*_*barrier*_*all* and *shmem*_*barrier* routines. See Section 8.5 and 8.5.1.
- Clarification about the completion and ordering semantics of *shmem*_*quiet* and *shmem*_*fence*. See Section 8.7.1 and 8.7.
- Clarifications about completion semantics of RMA and AMO routines. See Sections 8.3 and 8.4
- Clarifications on the memory model and the memory alignment requirements for symmetric data objects. See Section 3.
- Clarification on the execution model and the definition of a PE. See Section 4
- Clarifications of the semantics of *shmem*_*pe*_*accessible* and *shmem*_*addr*_*accessible*. See Section [8.1.3](#page-11-1) and [8.1.4.](#page-11-2)
- Added an annex on interoperability with MPI. See Annex [D.](#page-65-0)
- Added examples to the different interfaces.
- Clarification on the naming conventions for constant in *C* and *Fortran*. See Section [6](#page-7-6) and [8.6.1.](#page-49-1)
- Added API calls: *shmem*_*char*_*p*, *shmem*_*char*_*g*.
- Removed API calls: *shmem*_*char*_*put*, *shmem*_*char*_*get*.

1. VERSION 1.1 65

• The usage of *ptrdiff_t*, *size_t*, and *int* in the interface signature was made consistent with the description in Sections [8.5](#page-35-0) [8.3.3](#page-21-0) [8.3.6](#page-26-0)

PRAFT.