

Moving targets: Quantum circuit synthesis for fault-tolerant gate sets

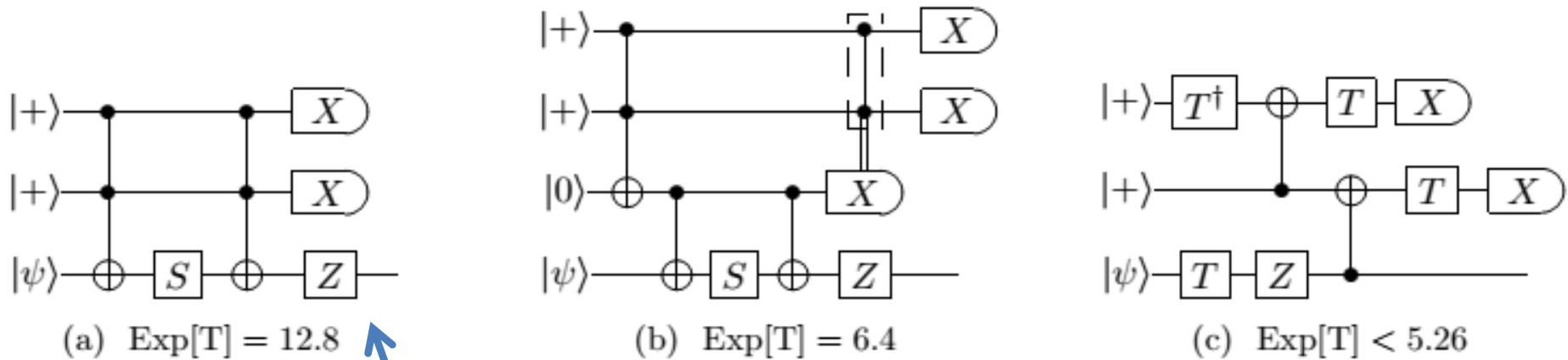
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Based on joint work with Alex Bocharov, Krysta Svore, and Nathan Wiebe
[arXiv:1404.5320](https://arxiv.org/abs/1404.5320) [arXiv:1406.2040](https://arxiv.org/abs/1406.2040) [arXiv:1409.3552](https://arxiv.org/abs/1409.3552)

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“Repeat-Until-Success” protocols



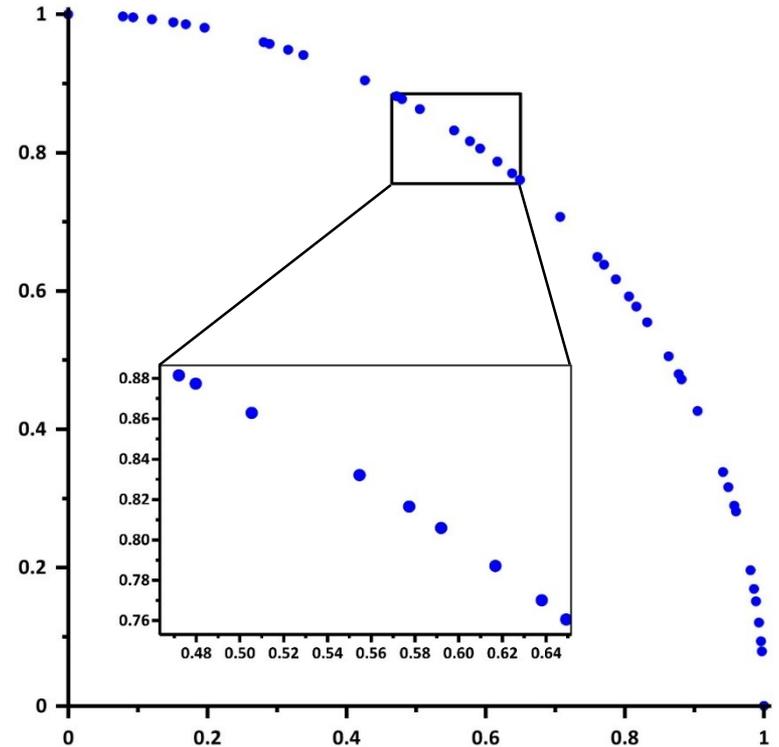
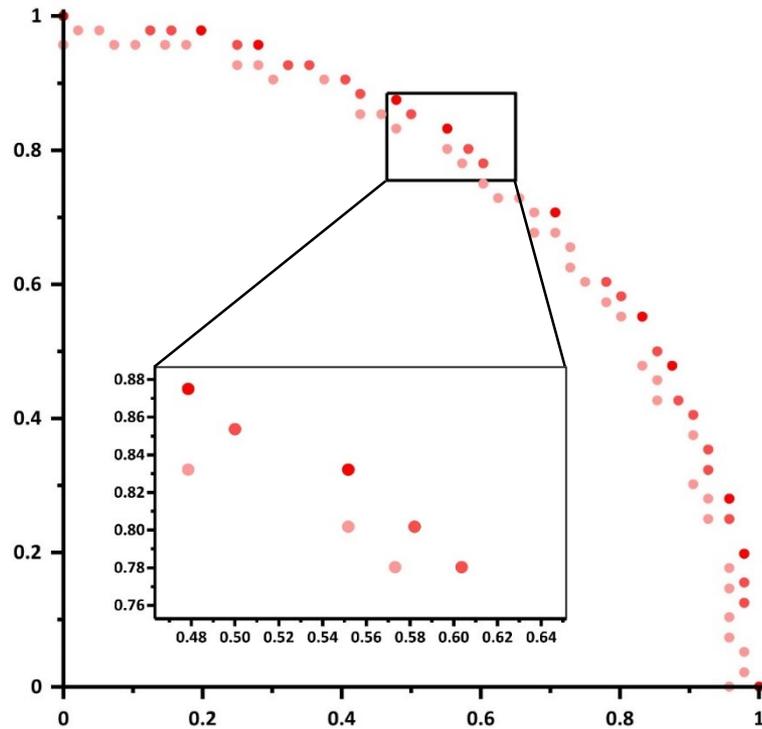
- Implements $V_3 = (I + 2iZ)/\sqrt{5}$ which corresponds to a z-rotation $R_z(\theta)$ with irrational angle θ .
- Probability of success $p_{\text{success}} = \frac{5}{8}$
- Failure branches: all lead to identity I_2 [Nielsen/Chuang'00]

Improved T-count by replacing Toffoli [Jones'12]

Found by enumeration [Paetznic/Svore'13]

RUS synthesis?

Approximations: unitaries versus RUS



a) Complete enumeration $\langle H, T \rangle$ unitaries of T-count $L \leq 8$ (i.e., $\ell \leq 4$). Red/grey denotes if within 0.98 of a z-rotation.

b) Enumeration of RUS protocols with expected T-count $L_{exp} \leq 7.5$. Blue points have $p_{success} \geq 0.8$

Cost advantage: $3.0 \log\left(\frac{1}{\epsilon}\right)$ for unitaries vs $1.15 \log\left(\frac{1}{\epsilon}\right)$ for RUS designs

Compiling z-rotations: overview

$$\begin{bmatrix} 1 & 0 \\ 0 & e^{i\theta} \end{bmatrix}$$

↓

$$\begin{bmatrix} z & 0 \\ 0 & z^* \end{bmatrix}$$

↓

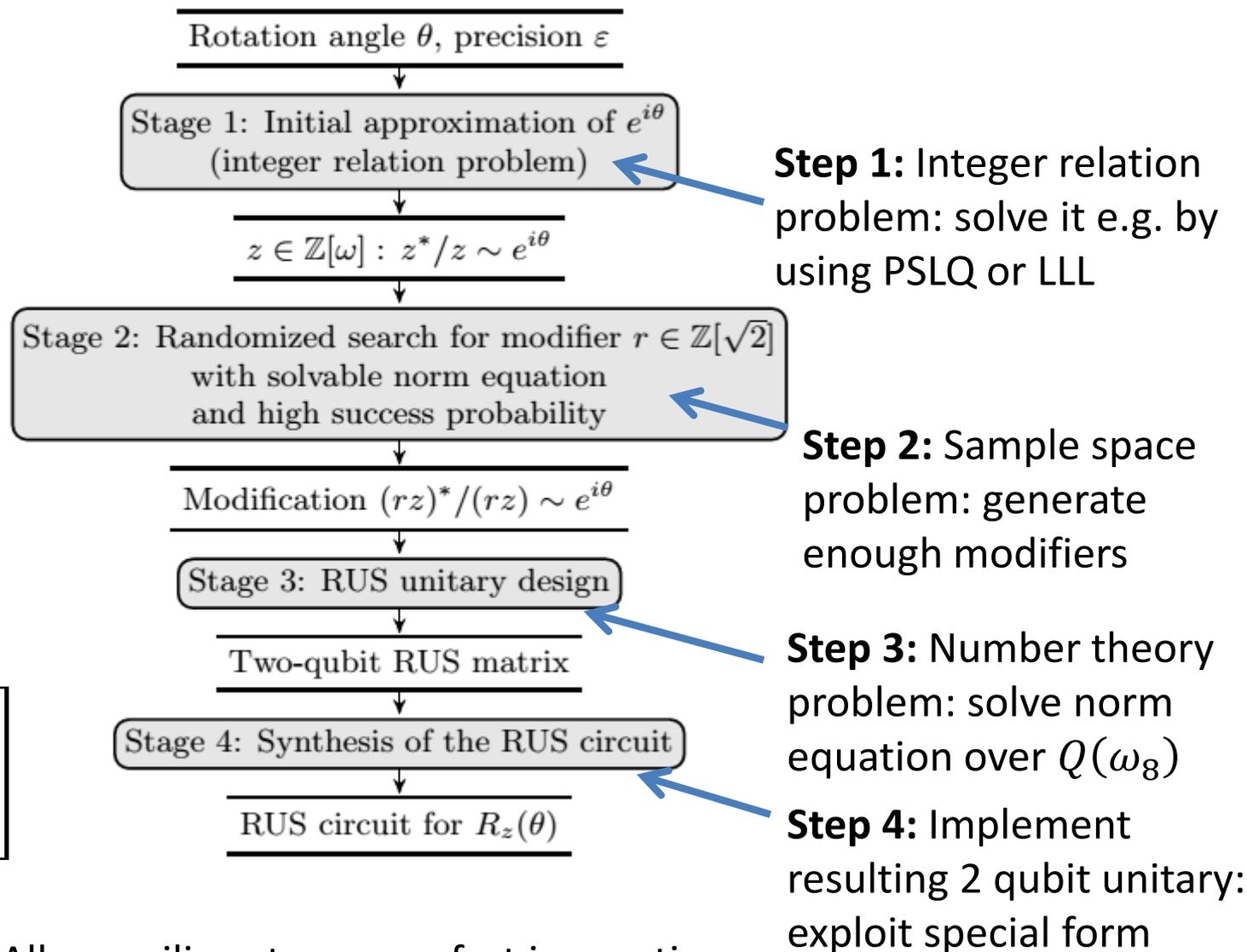
$$u = r z / (\sqrt{2})^L$$

↓

$$v = y / (\sqrt{2})^L$$

↓

$$\begin{bmatrix} u & -v^* & 0 & 0 \\ v & u^* & 0 & 0 \\ 0 & 0 & u^* & v^* \\ 0 & 0 & -v & u \end{bmatrix}$$



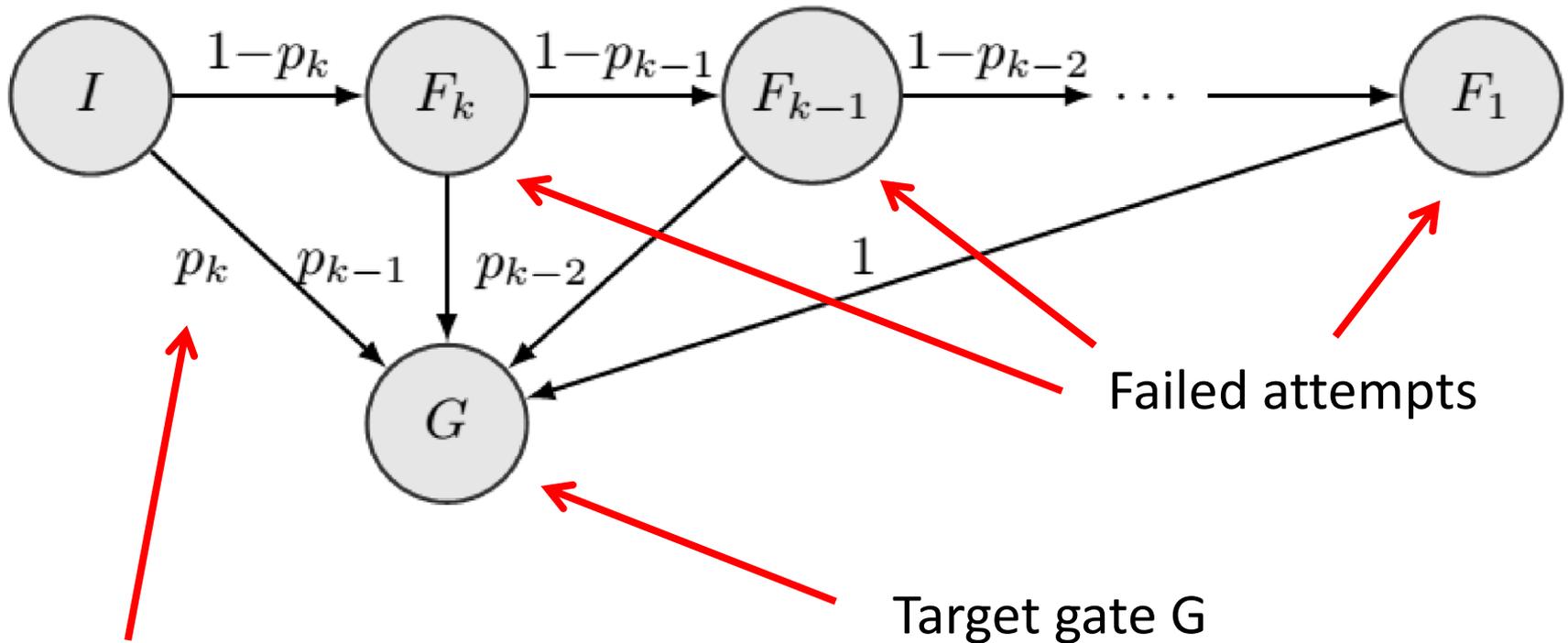
Compilation time: All compiling stages run fast in practice, even for extremely high accuracy (e.g., $\epsilon = 10^{-100}$).

Some open problems

- What are common features of the underlying fault-tolerant gate sets that can be leveraged for the development of an---ideally, systematic---approach toward quantum circuit synthesis?
- How can the synthesis layer inform the fault-tolerance layer about which gate sets are preferable over others?
- How can the trade space between resulting quantum circuit depth (or size) versus number of qubits versus error probability be explored in a meaningful way? Does it make sense to trade compilation time for quality of the resulting circuits?
- Does offline preparation help to bring down the complexity and if so, how can the resources be quantified in a meaningful way?
- How can the problem of synthesizing transformations in the multi-qubit case be addressed?

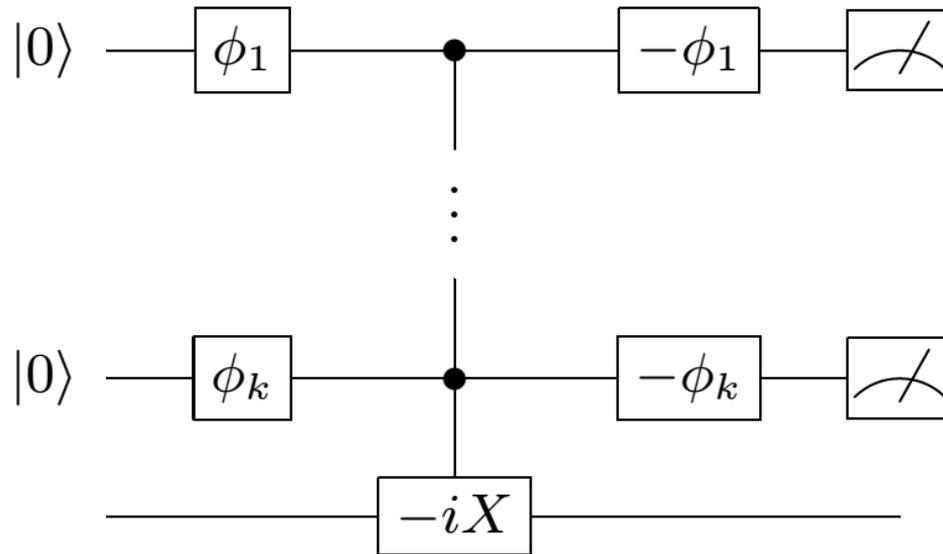
Backup slides

PQF synthesis = finding Markov chains



Success probabilities. Note that in “fallback” designs the last step has $P(\text{success})=1$

Arithmetic implemented RUS-style



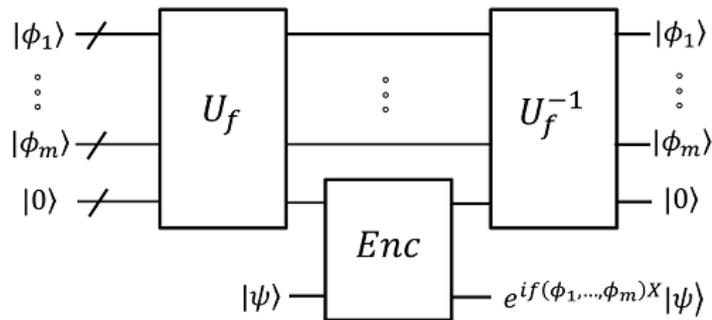
Gearbox circuit with k inputs, where gate ϕ_j denotes $e^{-i\phi_j X}$. Success is achieved if *every* measurement reads 0.

Fact: This defines an RUS circuit. Upon success it implements

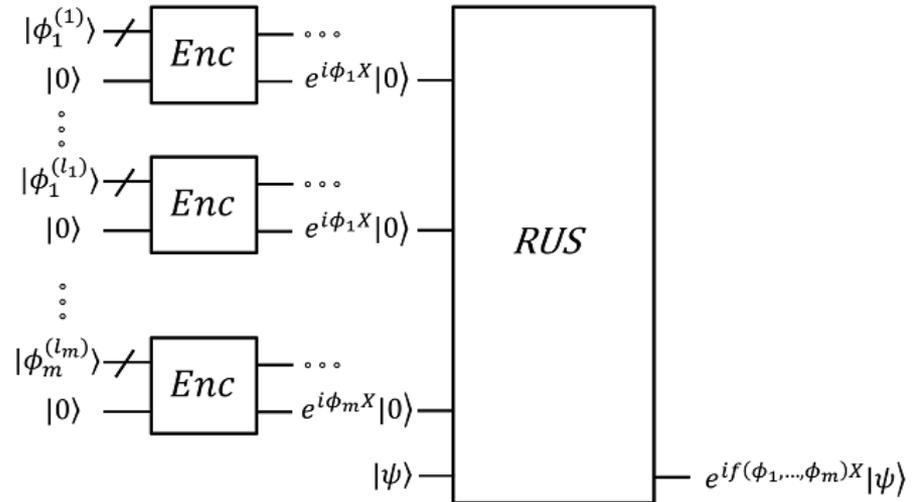
$$\text{GB} \left(e^{-i\phi_1 X} |0\rangle \cdots e^{-i\phi_k X} |0\rangle |\psi\rangle \right) = e^{-i \tan^{-1}(\tan^2(\arcsin(|\sin(\phi_1)| \cdots |\sin(\phi_k)|))} X |\psi\rangle .$$

Comparison: classical vs quantum

Classical/reversible:



Quantum using RUS synthesis:



Multiplier method	$n = 2$		$n = 4$		$n = 8$		$n = 16$	
	T -count	qubits						
Carry-ripple	2.34E+02	4	7.84E+02	8	2.80E+03	16	1.06E+04	32
Table-lookup	3.38E+03	2	3.26E+06	2	3.98E+09	2	1.13E+13	2
M_4	6.11E+01	3	1.97E+03	4	4.64E+04	4	3.00E+07	4
M_6	7.71E+02	4	1.67E+03	4	3.82E+03	4	5.21E+05	5